Domain Engineering A Basis for Safety Critical Software

Dines Bjørner Fredsvej 11, DK–2840 Holte, Denmark March 18, 2014: 08:22 1

• A Software Development Triptych:

Before software can be designed
we must have a reasonable grasp
of the requirements
that the software is supposed to fulfil.

And before requirements can be prescribed
we must have a reasonable grasp
of the "underlying" application domain.

• A Dogma:

- « Domain engineering now becomes
 - a software engineering development phase
 - ∞ in which a precise description,
 - ∞ desirably formal,
 - ∞ of the domain
 - ∞ within which the target software is to be embedded.
- Requirements engineering then becomes a phase of software engineering
 - ∞ in which one systematically derives requirements prescriptions ∞ from the domain description.

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- We illustrate the first element, \mathcal{D} , of this triptych $(\mathcal{D}, \mathcal{R}, \mathcal{S})$ by an example in which we show a description of a pipeline domain where, for example, the operations of pumps and valves are safety critical.
- We then summarise the methodological stages and steps of domain engineering.
- We finally weave considerations of *system safety criticality* into a section on domain facets.

- We believe this aspect of safety criticality is new:
 - \otimes We here connect safety criticality to domain engineering.
 - \otimes The study presented here need be deepened.
 - \otimes Similar connections need be made to
 - requirements engineering such as it can be "derived" from domain engineering, and to
 - ∞ the related software design.
 - ∞ That is, three distinct "layers" of safety engineering.

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1. Introduction

1.

• A Software Development Triptych:

Before software can be designed
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of the "underlying" application domain.

- **Domain engineering** now becomes
 - a software engineering development phase
 - « in which a precise description,
 - « desirably formal,
 - \otimes of the domain
 - \otimes within which the target software is to be embedded.

• **Requirements engineering** then becomes

- a phase of software engineering
- \otimes in which one systematically derives
 - ∞ requirements prescriptions
 - ∞ from the domain description —
 - ∞ carving out and extending, as it were, a subset of those
 - * domain properties that are computable and
 - * for which computing support is required.

• **Software design** is then

∞ the software engineering phase

« which results in code (and further documentation).

- We shall first
 - \otimes give a fairly large example, approximately 30 Slides,
 - « of a postulated domain of (say, oil or gas) pipelines;
 - \circledast the focus will be on ${\bf endurant}{\bf s}{:}$
 - ${\tt ∞}$ the observable **entities** that endure,
 - their **mereology**, that is, how they relate, and
 - their **attribute**s.
 - *** Perdurants: actions, events** and **behaviour**s will be very briefly mentioned.

- We shall then
 - \otimes on the background of this substantial example,
 - \otimes outline the basical principles, techniques and tools
 - \otimes for describing domains —
 - « focusig only on endurants.

- We shall review notions of **safety criticality**:
 - safety,
 failure,
 fault,
 risk.
- Other notions will also be briefly characterised:

 © component and
 © system

 safety, and
 - **stake-holder**,**∞ machine** and

« requirements.

- And, finally we shall detail the notion of **domain facet**s.
 - \otimes The various domain facets
 - ∞ somehow reflect domain views —
 - ∞ of logical or algebraic nature —
 - ∞ views that are shared across stake-holder groups,
 - ∞ but are otherwise clearly separable.
 - \otimes It is in connection with
 - the summary explanation of respective domain facets
 - that we identify respective **faults** and **hazards**.
 - \otimes The presentation is brief.

• We consider the following ideas new:

 \otimes the idea of describing domains before prescribing requirements \otimes and the idea of enumerating faults and hazards

 ∞ as related to individual facets.

 \otimes For the latter "discovery" we thank the organisers of ASSC 2014, notably Prof. Clive Victor Boughton.

2. An Example

- Our example is an abstraction of pipeline system endurants.
 - \otimes The presentation of the example
 - reflects a rigorous use of the domain analysis & description method outlined in Sect. 3,
 - but is relaxed with respect to not showing all one could say intermediate – analysis steps and description texts,
 - \ast but following stoichiometry ideas from chemistry

* makes a few short-cuts here and there.

- ∞ The use of the "stoichiometrical" reductions,
 - * usually skipping intermediate endurant sorts,
 - \ast ought properly be justified in each step —
 - \ast and such is adviced in proper, industry-scale analyses & descriptions.

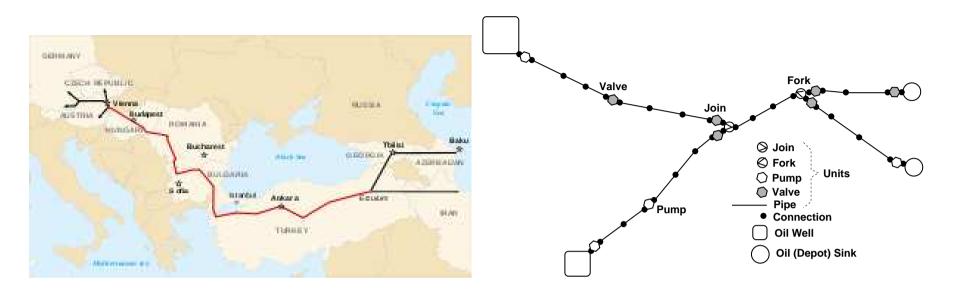


Figure 1: Pipelines. Flow is right-to-left in left figure, but left-to-right in right figure.

• The description only covers a few aspects of endurants.



Figure 2: Some oil pipeline system units: pump, pipe, valve

2.1. **Parts**

- 1. A pipeline system contains a set of pipeline units and a pipeline system monitor.
- 2. The well-formedness of a pipeline system depends on its mereology and the routing of its pipes.
- 3. A pipeline unit is either a well, a pipe, a pump, a valve, a fork, a join, or a sink unit.
- 4. We consider all these units to be distinguishable, i.e., the set of wells, the set pipe, etc., the set of sinks, to be disjoint.

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type

- 1. $\overline{PLS'}$, US, U, M
- 2. $PLS = \{ | pls:PLS' \cdot wf_PLS(pls) | \}$

<u>value</u>

- 2. wf_PLS: $PLS \rightarrow \underline{Bool}$
- 2. wf_PLS(pls) \equiv wf_Mereology(pls) \land wf_Routes(pls)
- 1. obs_Us: $PS \rightarrow U$ _set
- 1. $obs_M: PLS \rightarrow M$

type

- 3. $\overline{U} = We | Pi | Pu | Va | Fo | Jo | Si$
- 4. We :: Well
- 4. Pi :: Pipe
- 4. Pu :: Pump
- 4. Va :: Valv
- 4. Fo :: Fork
- 4. Jo :: Join
- 4. Si :: Sink

2.2. Part Identification and Mereology 2.2.1. Unique Identification

5. Each pipeline unit is uniquely distinguished by its unique unit identifier.

type

5. UI

<u>value</u>

5. $uid_UI: U \rightarrow UI$

axiom

5. $\forall pls:PLS, u, u': U \cdot \{u, u'\} \subseteq obs_Us(pls) \Rightarrow u \neq u' \Rightarrow uid_UI(u) \neq uid_UI(u')$

 $^{^{-1}}$ uid_UI is the unique identifier observer function for parts u:U. It is total. uid_UI(u) yields the unique identifier of u.

 $^{^{0}}$ The axiom expresses that for all pipeline systems all two distinct units, u, u' of such pipeline systems have distinct unique identifiers.

2.2.2. Unique Identifiers

6. From a pipeline system one can observe the set of all unique unit identifiers.

value

- 6. $xtr_UIs: PLS \rightarrow UI_set$
- 6. $xtr_UIs(pls) \equiv {uid_UI(u)|u:U \in obs_Us(pls)}$
- 7. We can prove that the number of unique unit identifiers of a pipeline system equals that of the units of that system.

theorem:

7. $\forall \text{ pls:PLS} \cdot \underline{\mathbf{card}} \text{ obs} \cdot Us(\text{pl}) = \underline{\mathbf{card}} \text{ xtr} \cdot UIs(\text{pls})$

 $^{^{0}}$ xtr_Uls is a total function. It extracts all unique unit identifiers of a pipeline system.

2.2.3. Mereology

- 8. Each unit is connected to zero, one or two other existing (formula line 8x.) input units and zero, one or two other existing (formula line 8x.) output units as follows:
 - a. A well unit is connected to exactly one output unit (and, hence, has no "input").
 - b. A pipe unit is connected to exactly one input unit and one output unit.
 - c. A pump unit is connected to exactly one input unit and one output unit.
 - d. A valve is connected to exactly one input unit and one output unit.
 - e. A fork is connected to exactly one input unit and two distinct output units.
 - f. A join is connected to exactly two distinct input units and one output unit.
 - g. A sink is connected to exactly one input unit (and, hence, has no "output").

| type | | |
|--|--|--|
| 8. $MER = UI - set \times UI - set$ | | |
| value | | |
| 8. mereo_U: $U \rightarrow MER$ | | |
| axiom | | |
| 8. wf_Mereology: $PLS \rightarrow \underline{Bool}$ | | |
| 8. wf_Mereology(pls) \equiv | | |
| 8. $\forall u: U \cdot u \in obs_Us(pls) \Rightarrow$ | | |
| 8x. <u>let</u> (iuis,ouis) = mereo_U(u) <u>in</u> iuis \cup ouis \subseteq xtr_UIs(pls) \wedge | | |
| 8. $\underline{case} (u, (\underline{card} \text{ iuis}, \underline{card} \text{ ouis})) \underline{of}$ | | |
| 8a. $(mk_We(we), (0,1)) \rightarrow \underline{true},$ | | |
| 8b. $(mk_Pi(pi),(1,1)) \rightarrow \underline{true},$ | | |
| 8c. $(mk_Pu(pu), (1,1)) \rightarrow \underline{true},$ | | |
| 8d. $(mk_Va(va),(1,1)) \rightarrow \underline{true},$ | | |
| 8e. $(mk_Fo(fo),(1,2)) \rightarrow \underline{true},$ | | |
| 8f. $(mk_Jo(jo), (2,1)) \rightarrow \underline{true},$ | | |
| 8g. $(mk_Si(si),(1,0)) \rightarrow \underline{true},$ | | |
| 8. $\underline{} \rightarrow \underline{\mathbf{false}} \ \underline{\mathbf{end}} \ \underline{\mathbf{end}}$ | | |

2.3. Part Concepts

- An aspect of domain analysis & description that was not covered in Sect. 2 was that of derived concepts.
- Example pipeline concepts are
 - \otimes routes,
 - \otimes acyclic or cyclic,
 - \otimes circular,

etcetera.

- In expressing well-formedness of pipeline systems
- one often has to develop subsidiary concepts such as these
- by means of which well-formedness is then expressed.

2.3.1. Pipe Routes

- 9. A route (of a pipeline system) is a sequence of connected units (of the pipeline system).
- 10. A route descriptor is a sequence of unit identifiers and the connected units of a route (of a pipeline system).

9. $\mathbf{R}' = \mathbf{U}^{\omega}$ 9. $\mathbf{R} = \{ | \mathbf{r}: \mathbf{Route'} \cdot \mathbf{wf}_{\mathbf{Route}}(\mathbf{r}) | \}$ 10. $\mathbf{RD} = \mathbf{UI}^{\omega}$ **axiom** 10. $\forall \mathbf{rd}: \mathbf{RD} \cdot \exists \mathbf{r}: \mathbf{R} \cdot \mathbf{rd} = \mathrm{descriptor}(\mathbf{r})$ **value** 10. $\mathrm{descriptor}: \mathbf{R} \to \mathbf{RD}$ 10. $\mathrm{descriptor}(\mathbf{r}) \equiv \langle \mathrm{uid}_{\mathbf{UI}}(\mathbf{r}[\mathbf{i}]) | \mathbf{i}: \mathbf{Nat} \cdot \mathbf{1} \leq \mathbf{i} \leq \mathbf{len} \mathbf{r} \rangle$ 11. Two units are adjacent if the output unit identifiers of one shares a unique unit identifier with the input identifiers of the other.

<u>value</u>

- 11. adjacent: $U \times U \rightarrow \underline{Bool}$
- 11. $adjacent(u,u') \equiv$
- 11. <u>**let**</u> (,ouis)=mereo_U(u),(iuis,)=mereo_U(u') <u>**in**</u>
- 11. ouis \cap iuis \neq {} <u>end</u>

- 12. Given a pipeline system, *pls*, one can identify the (possibly infinite) set of (possibly infinite) routes of that pipeline system.
 - a. The empty sequence, $\langle \rangle$, is a route of *pls*.
 - b. Let u, u' be any units of *pls*, such that an output unit identifier of u is the same as an input unit identifier of u' then $\langle u, u' \rangle$ is a route of *pls*.
 - c. If r and r' are routes of pls such that the last element of r is the same as the first element of r', then $r \hat{t} r'$ is a route of *pls*.
 - d. No sequence of units is a route unless it follows from a finite (or an infinite) number of applications of the basis and induction clauses of Items 12a.-12c.

value

- 12. Routes: $PLS \rightarrow RD$ -infset
- $Routes(pls) \equiv$ 12.
- 12a..
- let $rs = \langle \rangle \cup$ $\{\langle uid_UI(u), uid_UI(u') \rangle | u, u': U \cdot \{u, u'\} \subseteq obs_Us(pls) \land adjacent(u, u')\}$ 12b..
- \cup {r**tl** r'|r,r':R•{r,r'}⊆rs} 12c.
- 12d.. in rs end

2.3.2. Well-formed Routes

13. A route is acyclic if no two route positions reveal the same unique unit identifier.

<u>value</u>

- 13. acyclic_Route: $R \rightarrow \underline{Bool}$
- 13. $acyclic_Route(r) \equiv \sim \exists i,j: \underline{Nat} \cdot \{i,j\} \subseteq \underline{inds} r \land i \neq j \land r[i] = r[j]$

14. A pipeline system is well-formed if none of its routes are circular (and all of its routes embedded in well-to-sink routes).

<u>value</u>

- 14. wf_Routes: $PLS \rightarrow \underline{Bool}$
- 14. wf_Routes(pls) \equiv
- 14. non_circular(pls) \land are_embedded_in_well_to_sink_Routes(pls)
- 14. non_circular_PLS: $PLS \rightarrow \underline{Bool}$
- 14. non_circular_PLS(pls) \equiv
- 14. $\forall r: R \cdot r \in routes(p) \land acyclic_Route(r)$

15. We define well-formedness in terms of well-to-sink routes, i.e., routes which start with a well unit and end with a sink unit.

value

- 15. well_to_sink_Routes: $PLS \rightarrow R$ -set
- 15. well_to_sink_Routes(pls) \equiv
- 15. <u>**let**</u> rs = Routes(pls) <u>**in**</u>
- 15. $\{r | r: R \cdot r \in rs \land is_We(r[1]) \land is_Si(r[\underline{len} r])\} \underline{end}$

- 16. A pipeline system is well-formed if all of its routes are embedded in well-to-sink routes.
- 16. are_embedded_in_well_to_sink_Routes: $PLS \rightarrow \underline{Bool}$
- 16. are_embedded_in_well_to_sink_Routes(pls) \equiv
- 16. <u>**let</u>** wsrs = well_to_sink_Routes(pls) <u>**in**</u></u>
- 16. $\forall r: R \cdot r \in \text{Routes}(\text{pls}) \Rightarrow$
- 16. $\exists r':R,i,j:\underline{Nat}$.
- 16. $r' \in wsrs$

16.
$$\wedge \{i,j\} \subseteq \underline{inds} r' \land i \leq j$$

16. $\wedge \mathbf{r} = \langle \mathbf{r}'[\mathbf{k}] | \mathbf{k}: \mathbf{\underline{Nat}} \cdot \mathbf{i} \leq \mathbf{k} \leq \mathbf{j} \rangle \mathbf{\underline{end}}$

2.3.3. Embedded Routes

17. For every route we can define the set of all its embedded routes.

<u>value</u>

- 17. embedded_Routes: $R \rightarrow R$ -set
- 17. embedded_Routes(r) \equiv
- 17. $\{ \langle r[k] | k: \underline{\mathbf{Nat}} \cdot i \leq k \leq j \rangle \mid i, j: \underline{\mathbf{Nat}} \cdot i \{i, j\} \subseteq \underline{\mathbf{inds}}(r) \land i \leq j \}$

2.3.4. **A Theorem**

18. The following theorem is conjectured:

a. the set of all routes (of the pipeline system)

- b. is the set of all well-to-sink routes (of a pipeline system) and
- c. all their embedded routes

theorem:

18. \forall pls:PLS · 18. <u>let</u> rs = Routes(pls), 18. wsrs = well_to_sink_Routes(pls) <u>in</u> 18a.. rs = 18b.. wsrs \cup 18c.. \cup {{r'|r':R · r' \in embedded_Routes(r")} | r":R · r" \in wsrs} 17. end

2.4. Materials

19. The only material of concern to pipelines is the gas¹ or liquid² which the pipes transport³.

type 19. GoL value 19. obs_GoL: U → GoL

¹Gaseous materials include: air, gas, etc. ²Liquid materials include water, oil, etc. ³The description of this document is relevant only to gas or oil pipelines.

2.5. Attributes 2.5.1. Part Attributes

20. These are some attribute types:

a. estimated current well capacity (barrels of oil, etc.),

b. pipe length,

- c. current pump height,
- d. current valve open/close status and
- e. flow (e.g., volume/second).

| type | |
|------|------------------------|
| 20a | WellCap |
| 20b | LEN |
| 20c | Height |
| 20d | ValSta == open close |
| 20e | Flow |

21. Flows can be added (also distributively) and subtracted, and22. flows can be compared.

value

- 21. \oplus, \ominus : Flow \times Flow \rightarrow Flow
- 21. \oplus : Flow<u>-set</u> \rightarrow Flow

22.
$$<,\leq,=,\neq,\geq,>$$
: Flow × Flow \rightarrow **Bool**

- 23. Properties of pipeline units include
 - a. estimated current well capacity (barrels of oil, etc.),
 - b. pipe length,
 - c. current pump height,
 - d. current valve open/close status,
 - e. current \mathcal{L} aminar in-flow at unit input,
 - f. current \mathcal{L} aminar in-flow leak at unit input,
 - g. maximum \mathcal{L} aminar guaranteed in-flow leak at unit input,
 - h. current \mathcal{L} aminar leak unit interior,
 - i. current \mathcal{L} aminar flow in unit interior,
 - j. maximum \mathcal{L} aminar guaranteed flow in unit interior,
 - k. current \mathcal{L} aminar out-flow at unit output,
 - l. current \mathcal{L} aminar out-flow leak at unit output,
 - m. maximum guaranteed $\mathcal L$ aminar out-flow leak at unit output.

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<u>value</u>

| 23a | attr_WellCap: We \rightarrow WellCap |
|-----|---|
| 23b | attr_LEN: $Pi \rightarrow LEN$ |
| 23c | attr_Height: $Pu \rightarrow Height$ |
| 23d | attr_ValSta: Va \rightarrow VaSta |
| 23e | $\operatorname{attr_In_Flow}_{\mathcal{L}}: U \to UI \to Flow$ |
| 23f | $\operatorname{attr_In_Leak}_{\mathcal{L}}: U \to UI \to Flow$ |
| 23g | $\operatorname{attr}_{\operatorname{Max}}\operatorname{In}_{\operatorname{Leak}} \mathcal{L}: U \to UI \to Flow$ |
| 23h | $\operatorname{attr_body_Flow}_{\mathcal{L}}: U \to \operatorname{Flow}$ |
| 23i | $\operatorname{attr_body_Leak}_{\mathcal{L}}: U \to Flow$ |
| 23j | $\operatorname{attr}_{\operatorname{Max}}\operatorname{Flow}_{\mathcal{L}}: U \to \operatorname{Flow}$ |
| 23k | $\operatorname{attr_Out_Flow}_{\mathcal{L}}: U \to UI \to Flow$ |
| 23l | $\operatorname{attr_Out_Leak}_{\mathcal{L}}: U \to UI \to Flow$ |
| 23m | $\operatorname{attr}_{\operatorname{Max}}\operatorname{Out}_{\operatorname{Leak}}\mathcal{L}: \operatorname{U} \to \operatorname{UI} \to \operatorname{Flow}$ |

2.5.2. Flow Laws

24. "What flows in, flows out !". For *L*aminar flows: for any non-well and non-sink unit the sums of input leaks and in-flows equals the sums of unit and output leaks and out-flows.

Law:

- 24. \forall u:U\We\Si ·
- 24. $sum_in_leaks(u) \oplus sum_in_flows(u) =$
- 24. $\operatorname{attr_body_Leak}_{\mathcal{L}}(u) \oplus$
- 24. $sum_out_leaks(u) \oplus sum_out_flows(u)$

<u>value</u>

```
sum_in_leaks: U \rightarrow Flow
sum_in_leaks(u) \equiv
     let (iuis,) = mereo_U(u) in
     \oplus \{ attr_In_Leak_{\mathcal{L}}(u)(ui) | ui: UI \cdot ui \in iuis \} <u>end</u>
sum_in_flows: U \rightarrow Flow
sum_in_flows(u) \equiv
     let (iuis,) = mereo_U(u) in
     \oplus \{ attr_In_Flow_{\mathcal{L}}(u)(ui) | ui: UI \cdot ui \in iuis \} <u>end</u>
sum_out_leaks: U \rightarrow Flow
sum_out_leaks(u) \equiv
     \underline{\mathbf{let}} (,ouis) = mereo_U(u) \underline{\mathbf{in}}
     \oplus  {attr_Out_Leak<sub>L</sub>(u)(ui)|ui:UI·ui \in ouis} end
sum_out_flows: U \rightarrow Flow
sum_out_flows(u) \equiv
     \underline{\mathbf{let}} (,ouis) = mereo_U(u) \underline{\mathbf{in}}
     \oplus  {attr_Out_Leak<sub>\mathcal{L}</sub>(u)(ui)|ui:UI \cdot ui \in ouis} end</sub>
```

25. "What flows out, flows in !". For *L*aminar flows: for any adjacent pairs of units the output flow at one unit connection equals the sum of adjacent unit leak and in-flow at that connection.

Law:

- 25. $\forall u, u': U \cdot adjacent(u, u') \Rightarrow$
- 25. <u>**let**</u> (,ouis)=mereo_U(u), (iuis',)=mereo_U(u') <u>**in**</u>
- 25. <u>assert</u>: $uid_U(u') \in ouis \land uid_U(u) \in iuis '$
- 25. $\operatorname{attr_Out_Flow}_{\mathcal{L}}(u)(\operatorname{uid}_U(u')) =$
- 25. $\operatorname{attr_In_Leak}_{\mathcal{L}}(u)(\operatorname{uid_U}(u)) \oplus \operatorname{attr_In_Flow}_{\mathcal{L}}(u')(\operatorname{uid_U}(u)) \underline{end}$

2.5.3. Open Routes

26. A route, r, is open

a. if all values, v, of the route are open and

b. if all pumps, p, of the route are pumping.

value

- 26. is_open: $R \rightarrow \underline{Bool}$
- 26. is_open(r) \equiv
- 26a.. $\forall mkPu(p):Pu \cdot mkPu(p) \in \underline{elems} r \Rightarrow is_pumping(p) \land$
- 26b.. $\forall mkVa(v): Va \cdot mkVa(v) \in \underline{elems} r \Rightarrow is_open(v)$

2.6. Domain Perdurants

2.6.1. **Actions**

- We shall not formalise any specific actions.
- Informal examples of actions are:
 - \otimes opening and closing a well,
 - \otimes start and stop pumping,
 - \otimes open and close values,
 - \otimes opening and closing a sink and
 - \otimes sense current unit flow.

2.6.2. **Events**

- We shall not formalise any specific events.
- Informal examples of events are:
 - \otimes empty well,
 - \otimes full sink,
 - \otimes start pumping signal to pump with no liquid material,
 - « pump ignores start/stop pumping signal,
 - ∞ valve ignores opening/closing signal,
 - \otimes excessive to catastrophic unit leak, and
 - \otimes unit fire or explosion.

2.6.3. Behaviours

- We shall not formalise any specific behaviours.
- Informal examples of behaviours are:
 - \otimes start pumping and opening up values across a pipeline system, and
 - \otimes stop pumping and closing down values across a pipeline system.

3. Basic Domain Description

- In this section and the next we shall survey basic principles of describing, respectively,
 - \otimes domain intrinsics and other
 - \otimes domain facets.

- By an **entity** we shall understand a **phenomenon**
 - ∞ that can be **observe**d, i.e., be

© seen or

© touched

by humans,

- \otimes or that can be ${\sf conceived}$
 - ∞ as an abstraction
 - ∞ of an entity.

- The method can thus be said to provide the *domain analysis prompt*:
 - $\otimes is_entity$
 - \otimes where is_entity(θ) holds if θ is an entity.

• A **domain** is characterised by its

 \otimes observable, i.e., manifest entities \otimes and their qualities .

• By a **quality** of an entity we shall understand

- ∞ a **property** that can be given a *name* and
- \otimes whose value can be
 - precisely measured by physical instruments
 or otherwise identified.
- Example: Unique identifiers (Slide 20, Item 5.), mereology (Slide 22, Item 8.) and the well capacity (Slide 35, Item 20a..), pipe length (Slide 35, Item 20b..), current pump height (Slide 35, Item 20c..), current valve open/close status (Slide 35, Item 20d..) and flow (Slide 35, Item 20e..) attributes are qualities ■

- By a sort (or type) we shall understand
 & the largest set of entities
 & all of which have the same qualities.
- By an endurant entity (or just, an endurant) we shall understand
 anything that can be observed or conceived,
 as a "complete thing",
 at no matter which given snapshot of time.
- Thus the method provides a *domain analysis prompt*:
 - $\circledast \texttt{is_endurant}$ where
 - \otimes is_endurant(e) holds if entity e is an endurant.

- By a **perdurant entity** (or just, an perdurant) we shall understand
 - \otimes an entity
 - « for which only a fragment exists
 - if we look at or touch them
 - at any given snapshot in time, that is,
 - were we to freeze time we would only see or touch
 a fragment of the perdurant.
- Thus the method provides a *domain analysis prompt*:
 - **wis_perdurant** where
 - $\otimes is_perdurant(e)$ holds if entity e is a perdurant.

By a discrete endurant we shall understand something which is
* separate or distinct in form or concept,
* consisting of distinct or separate parts.

• Thus the method provides a *domain analysis prompt*:

 $\circledast \texttt{is_discrete}$ where

 \otimes is_discrete(e) holds if entity e is discrete.

• By a continuous endurant

 \otimes we shall understand something which is

- « prolonged without interruption,
- \otimes in an unbroken series or pattern.

We use the term **material** for continuous endurants.

• Thus the method provides a *domain analysis prompt*:

*** is_continuous** where

 \otimes is_continuous(e) holds if entity e is a continuous entity.

3.0.1. Endurant Entities

We distinguish between endurant and perdurant entities. 3.0.1.1 Parts and Materials

- The manifest entities, i.e., the endurants, are called
 - \otimes parts, respectively

 \otimes materials.

• We use the term **part** for discrete endurants,

 $\text{ w that is: } is_part(p) \equiv is_endurant(p) \land is_discrete(p).$

 \bullet We use the term material for continuous endurants.

- Discrete endurants are
 - « either atomic
 - ⇔ or composite.
- By an **atomic endurant** we shall understand
 - \otimes a discrete endurant which
 - « in a given context,
- The method can thus be said to provide the *domain analysis prompt*:
 - \otimes is_atomic where is_atomic(p) holds if p is an atomic part.
- **Example**: Pipeline units, U, and the monitor, M, are considered atomic

• By a **composite endurant** we shall understand

- \otimes a discrete endurant which
- \otimes in a given context,
- The method can thus be said to provide the *domain analysis prompt*:
 - \otimes is_composite where is_composite(p) holds if p is an a composite part.
- **Example**: The pipeline system, PLS, and the set, Us, of pipeline units are considered composite entities

3.0.1.2 Part Observers

- From atomic parts we cannot observe any sub-parts.
- But from composite parts we can.
- For composite parts, p, the *domain description prompt*
 observe_part_sorts(p)
- yields some *formal description text* according to the following *schema*:

<u>type</u> P₁, P₂, ..., P_n;⁴ <u>value</u> obs_P₁: P \rightarrow P₁, obs_P₂: P \rightarrow P₂,...,obs_P_n: P \rightarrow P_n;⁵

⁵Thus RSL <u>value</u> clause defines n function values. All from type P into some type P_i .

⁴This RSL <u>type</u> clause defines P_1 , P_2 , ..., P_n to be sorts.

• where

- \otimes sort names P₁, P₂, ..., P_n
- ∞ are chosen by the domain analyser,

∞ may have been defined

already,

∞ but not recursively

- ∞ must denote disjoint sorts, and be discharged to secure disjointness of sorts.
- **Example**: Three formula lines (Slide 18, Items 1.) illustrate the basic sorts (PLS', US, U, M) and observers (obs_US, obs_M) of pipeline systems .

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3.0.1.3 Sort Models

- A part sort is an **abstract type**.
 - \otimes Some part sorts, P, may have a concrete type model, T.
 - « Here we consider only two such models:

 ∞ one model is as sets of parts of sort A: T = A-set;

- ∞ the other model has parts being of either of two or more alternative, disjoint sorts: T=P1|P2|...|PN.
- The *domain analysis prompt*:
 - \otimes has_concrete_type(p)
- \bullet holds if part p has a concrete type.
- In this case the *domain description prompt*

 $\otimes observe_concrete_type(p)$

𝔅 yields some *formal description text* according to the following *schema*,

* either

where $\mathcal{E}(...)$ is some type expression over part sorts and where P1,P2,...,PN are either (new) part sorts or are auxiliary (abstract or concrete) types⁸; * or:

```
    \underline{type} \\
    T = P1 | P2 | ... | PN<sup>9</sup> \\
    P_1, P_2, ..., P_n \\
    P1 :: mkP1(P_1), P2 :: mkP2(P_2), ..., PN :: mkPN(P_n) <sup>10</sup> 

    \underline{value} \\
    obs_T: P \to T^{11}
```

⁶The concrete type definition T = *E*(P1,P2,...,PN) define type T to be the set of elements of the type expressed by type expression *E*(P1,P2,...,PN).
⁷**obs**_T is a function from any element of P to some element of T.
⁸ The *domain analysis prompt*: sorts_of(t) yields a subset of {P1,P2,...,PN}.

 ${}^{\circ}A|B$ is the union type of types A and B.

¹⁰Type definition A :: mkA(B) defines type A to be the set of elements mkA(b) where b is any element of type B

¹¹**obs**_T is a function from any element of P to some element of T.

3.0.1.4 Material Observers

• Some parts p of sort P may contain material.

 $\circledast \ensuremath{\operatorname{The}}$ domain analysis prompt

material(p)

 \otimes holds if composite part p contains one or more materials.

- The *domain description prompt* observe_material_sorts(*p*)
- yields some *formal description text* according to the following *schema*:

<u>type</u> M₁, M₂, ..., M_m; value obs_M₁: $P \rightarrow M_1$, obs_M₂: $P \rightarrow M_2$, ..., obs_M_m: $P \rightarrow M_m$;

 \otimes where values, m_i , of type M_i satisfy is_material(m) for all i; \otimes and where $M_1, M_2, ..., M_m$ must be disjoint sorts.

• **Example**: We refer to Slide 34, Item 19.

3.0.2. Endurant Qualities

- We have already, above, treated the following properties of endurants:
 - $\otimes \texttt{is_discrete},$
 - $\otimes \texttt{is}_\texttt{continuous},$
 - $\otimes \texttt{is}_\texttt{atomic},$
 - $\circledast \texttt{is_composite}$ and
 - ∞ has_material.
- We may think of those properties as **external qualities**.
- In contrast we may consider the following internal qualities:
 - w has_unique_identifier (parts),
 - **whas_mereology** (parts) and
 - **whas_attributes** (parts and materials).

3.0.2.1 Unique Part Identifiers

- \bullet Without loss of generality we can assume that every part has a unique identifier $^{12}.$

 - \otimes If two parts are claimed to have the same unique identifier then they are identical.
- The domain description prompt: observe_unique_identifier(p)
- yields some *formal description text* according to the following *schema*:

```
\underline{\underline{\mathbf{type}}}_{\underline{\mathbf{value}}} \begin{array}{c} \mathrm{PI};\\ \underline{\mathbf{value}} \end{array} \mathbf{uid}_{-}\mathrm{P:} \ \mathrm{P} \rightarrow \mathrm{PI}; \end{array}
```

• **Example**: We refer to Slide 20, Item 5.

¹²That is, **has_unique_identifier(**p**)** for all parts p.

3.0.2.2 Part Mereology

• By **mereology** [Lesniewski1] we shall understand

- \otimes the study, knowledge and practice of
- \otimes parts,
- \otimes their relations to other parts
- \otimes and "the whole" $\mbox{.}$
- Part relations are such as:
 - « two or more parts being connected,
 - \otimes one part being embedded within another part, and
 - ∞ two or more parts sharing attributes.

• The *domain analysis prompt*:

has_mereology(p)

- holds if the part p is related to some others parts (p_a, p_b, \ldots, p_c) .
- The *domain description prompt*: observe_mereology(p) can then be invoked and
- yields some *formal description text* according to the following *schema*:

<u>type</u> MT = $\mathcal{E}(\text{PI}_A, \text{PI}_B, ..., \text{PI}_C);$ <u>value</u> mereo_P: P \rightarrow MT;

where $\mathcal{E}(...)$ is some type expression over unique identifier types of one or more part sorts.

- Mereologies are expressed in terms of structures of unique part identifiers.
- Usually mereologies are constrained. Constraints express
 - \otimes that a mereology's unique part identifiers must indeed reference existing parts, but also
 - \otimes that these mereology identifiers "define" a proper structuring of parts.
- **Example**: We refer to Items 8.–8g.. Slides 22–23

3.0.2.3 Part and Material Attributes

• Attributes are what really endows parts with qualities.

 \otimes The external properties¹³

- « are far from enough to distinguish one sort of parts from another.
- ∞ Similarly with unique identifiers and the mereology of parts.
- We therefore assume, without loss of generality, that
 - \otimes every part, whether discrete or continuous,
 - \otimes whether, when discrete, atomic or composite,

 \otimes has at least one attribute.

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- ∞ is_continuous,
- ∞ is_atomic,

is_composite
 whas_material.

• By an **endurant attribute**, we shall understand

- \otimes a property that is associated with an endurant e of sort E,
 - ∞ and if removed from endurant e,
 - ${\scriptstyle \scriptsize \ensuremath{\varpi}}$ that endurant would no longer be endurant e
 - ∞ (but may be an endurant of some other sort E'); and
- where that property itself has no physical extent (i.e., volume),a sthe endurant may have,
 - ∞ but may be measurable by physical means.

- The *domain description prompt* observe_attributes(p)
- yields some *formal description text* according to the following *schema*:

 $\underline{\mathbf{type}}_{\mathbf{value}} A_1, A_2, \dots, A_n, ATTR; \\ \underline{\mathbf{value}}_{\mathbf{value}} \mathbf{attr}_A_1: P \rightarrow A_1, \mathbf{attr}_A_2: P \rightarrow A_2, \dots, \mathbf{attr}_A_n: P \rightarrow A_n, \\ \mathbf{attr}_ATTR: P \rightarrow ATTR;$

• where <u>for</u> \forall p:P, attr_A_i(attr_ATTR(p)) \equiv attr_A_i(p).

• **Example**: We refer to Slides 35–38

3.0.3. Perdurant Entities

- But we shall summarise one essence for the description of perdurants.
- There is a notion of **state**.

 - \otimes Examples of such qualities are
 - ∞ the mereology of a part, and
 - ∞ part attributes whose value may change.

- There is the notion of **function signature**.
 - \otimes A function signature, f: A ($\rightarrow | \stackrel{\sim}{\rightarrow}$) R,
 - ∞ gives a name, say f, to a function,
 - ∞ expresses a type, say T_A , of the arguments of the function, ∞ expresses whether the function is total (\rightarrow) or partial $(\stackrel{\sim}{\rightarrow})$, and ∞ expresses a type, say T_R , of the result of the function.

- There is the notion of **channel**s of synchronisation & communication between behaviours.
 - \otimes Channels have names, e.g., ch, ch_i, ch_o.
 - $\label{eq:channel names appear in the signature of behaviour functions:} \underline{value} \ b: \ A \to \underline{in} \ ch_i \ \underline{out} \ ch_o \ R.$
 - $\lim ch_i$ indicates that behaviour **b** may express willingness to communicate an input message over channel ch_i ; and
 - \otimes <u>out</u> ch_o indicates that behaviour b may express an offer to communicate an output message over channel ch_o.

• There is a notion of **function pre/post-conditions**.

 \otimes A function pre-condition is a predicate over argument values.

Ø

Ø

 \odot

 \odot

 \otimes

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- Action signatures
 - \otimes include states, Σ , in both
 - \otimes arguments, $A \times \Sigma$, and results, Σ :
 - \ll f: A \times Σ \rightarrow Σ ;

 \otimes f denotes a function in the function space $A \times \Sigma \rightarrow \Sigma$.

• Action pre/post-conditions:

<u>value</u> $f(a,\sigma) \underline{as} \sigma'; \underline{pre}: \mathcal{P}_f(a,\sigma); \underline{post}: \mathcal{Q}_f(a,\sigma,\sigma')$

 \otimes have predicates \mathcal{P}_f and \mathcal{Q}_f \otimes delimit the value of f within that function space; $\otimes \mathcal{P}_f$ • Event signatures

 \otimes are typically predicates from pairs of before and after states: \otimes e: $\Sigma \times \Sigma \rightarrow \underline{Bool}$.

• Event pre/post-conditions

<u>value</u> e: $\Sigma \times \Sigma \rightarrow \underline{Bool}$; $e(\sigma, \sigma') \equiv \mathcal{P}_e(\sigma) \land \mathcal{Q}_e(\sigma, \sigma')$

- \otimes have predicates \mathcal{P}_e and \mathcal{Q}_e
- \otimes delimit the value of **e** within the $\Sigma \times \Sigma \rightarrow \underline{Bool}$ function space;
- $\otimes \mathcal{P}_e$ characterises states leading to event \mathbf{e} ;
- $\otimes \mathcal{Q}_e$ characterises states, σ' , resulting from the event caused by σ .

- In principle we can associate a behaviour with every part of a domain.
 - \otimes Parts, p, are characterised by their unique identifiers, pi:PI and a state, attrs:ATTRS.
 - \otimes We shall, with no loss of generality, assume part behaviours to be never-ending.
 - The unique part identifier, pi:Pl, and its the part mereology, say $\{ pi_1, pi_2, ..., pi_n \}$, determine a number of channels

 - \otimes able to communicate messages of $\mathbf{type}\ \mathsf{M}.$

• Behaviour signatures:

b: pi:PI × ATTR \rightarrow <u>in</u> in_chs <u>out</u> out_chs <u>Unit</u>

 \otimes then have

 ∞ input channel expressions in_chs and

- o output channel expressions out_chs
- ∞ be suitable predicates over
- $(chs[pi,pi_j]|j:\{1,2,...,n\} \}.$
- \otimes <u>Unit</u> designate that **b** denote a never-ending process.
- We omit dealing with behaviour pre-conditions and invariants.

4. Interlude

- We have covered one aspect of the modelling of one set of domain entities, the intrinsic facets of endurants.
 - « For the modelling of perdurants we refer to a forthcoming report.
- In the next section we shall survey the modelling of further domain facets.
- We shall accompany this survey to a survey of safety issues.
- To do so in a reasonably coherent way we need establish a few concepts:
 - \circledast the safety notions of

• failure, *• error* and *• fault*; *• the notion of stakeholder and • the notion of requirements.*

4.1. Safety-related Concepts

Some characterisations are:

4.1.0.1 Safety

By *safety*, in the context of a domain being dependable, we mean

• some measure of continuous delivery of service of

« either correct service,

- « or incorrect service after benign failure,
- that is: measure of time to catastrophic failure.

4.1.0.2 **Failure**

- A domain *failure* occurs
- when the delivered service
- deviates from fulfilling the domain function,
- the latter being what the domain is aimed at.

4.1.0.3 Error

• An *error*

- \bullet is that part of a domain state
- which is liable to lead to subsequent failure.
- An error affecting the service
- is an indication that a failure occurs or has occurred.

4.1.0.4 Fault

- The adjudged (i.e., the 'so-judged')
- \bullet or hypothesised cause of an error
- is a *fault*.

$4.1.0.5 \,\, \text{Hazard}$

• A hazard is

any source of potential damage, harm or adverse health effects on something or someone under certain conditions at work.

$4.1.0.6 \,\, {\rm Risk}$

• A **risk** is

- \otimes the chance or probability that a person
- \otimes will be harmed or experience an adverse health effect
- \otimes if exposed to a hazard.
- \otimes It may also apply to situations with property or equipment loss.

4.1.0.7 Faults and Hazards

- The concept of hazard is not the same as the concept of fault.
- *"System safety takes a larger view of hazards than just failures*¹⁴*:*
 - Hazards are not always caused by failures, and all failures do not cause hazards.
 - « Serious accidents have occurred
 - $\ensuremath{{\ensuremath{\scriptsize \odot}}}$ while system components
 - were all functioning exactly as specified,
 - [®] that is, without failure.
 - If failures only are considered in a safety analysis, many potential accidents will be missed.
 - In addition, the engineering approaches to preventing failures (increasing reliability) and preventing hazards (increasing safety) are different and sometimes conflict."

¹⁴Leveson: [White Paper on Approaches to Safety Engineering]

4.2. System and Component Safety

- There appears to be a number of safety concepts 15 :
 - \otimes component safety,
 - \otimes industrial safety,
 - \otimes reliability, and
 - \otimes system safety.
- We shall focus on component and system safety.

¹⁵Leveson: [White Paper on Approaches to Safety Engineering]

4.2.0.1 Component

• By a **component** we shall understand

- \otimes basically the same as an atomic part
- \otimes together with actions, events and behaviours
 - whose state is anchored
 - ∞ in one or more attributes of that part,
- \otimes such that these actions, etc.,
 - do nor involve other component or [sub]system states.
- \otimes That is, "componentry"

excludes considerations of shared attributes.

4.2.0.2 System

• By a **system** or **sub-system** we shall understand

- \otimes basically the same as a composite part
- \otimes together with actions, events and behaviours
 - ∞ whose state is anchored
 - ∞ in one or more attributes
 - * of that part
 - * as well as of one or more other parts.
- \otimes That is, "system-hood"

presumes considerations of shared attributes.

4.2.0.3 System Safety

- "The primary concern of system safety 16
 - *∞* is the management of hazards:
 - their identification,
 - evaluation,
 - \odot elimination, and
 - o control
 - through
 - o analysis,

• design and

management
 procedures."

¹⁶Leveson: [White Paper on Approaches to Safety Engineering]

- "System safety deals with systems as a whole rather than with subsystems or components ¹⁷:
 - « Safety is an emergent property of systems,
 - « not a component property.
 - « One of the principle responsibilities of system safety is
 - « to evaluate the interfaces between the system components
 - « and determine the effects of component interaction,
 - $\ensuremath{\circledast}$ where the set of components includes
 - humans,
 machines, and
 the environment."
- The system interfaces are given by the mereology.

¹⁷Leveson: [White Paper on Approaches to Safety Engineering]

4.2.0.4 Component Safety

- For a component,
- that is, an atomic part,
- we can, at most, speak of faults
- when considering safety.¹⁸

¹⁸The borderline between hazards that are not faults and faults is too vague.

4.3. Stake-holder

• By a **domain stake-holder** we shall understand

- **Examples**: The following are examples of pipeline stake-holders:
 - \otimes the owners of the pipeline,
 - \otimes the oil or gas companies using the pipeline,
 - ∞ the pipeline managers and workers,
 - « the owners and neighbours of the lands occupied by the pipeline,
 - \otimes the citizens possibly worried about gas- or oil pollution,
 - \otimes the state authorities regulating and overseeing pipelining,
 - etcetera

4.4. Machines and Requirements 4.4.1. Machine

• By the **machine** we shall understand

- \otimes the combination of
- hardware, say computers and communication, and software.

4.4.2. Requirements

- By a **requirements** we understand (cf. IEEE Standard 610.12):
 - \circledast "A condition or capability needed by a user to solve a problem or achieve an objective" .
- We shall think only of requirements as requirements to a machine.
- We can now "repeat" the definitions
 - « of safety, failure, error and fault given above,
 - \otimes but now with the term
 - ∞ 'domain' replaced by the term 'machine'
 - ∞ (sometimes with the term 'domain+machine').
- This then becomes the context in which most safety criticality is discussed.

- We shall not cover requirements in this talk.
- We refer to
 - [Bjørner: From Domains to Requirements, 2008].
 - ∞ That paper describes how to "derive"
 - \otimes systematically, but, of course, not automatically
 - \otimes major parts of requirements prescriptions from a domain descriptions.
- Thus we shall not cover the classical approach to safety analysis.

 - © One in which first get an as complete as possible overview of "all" safety aspects of a domain.

5. Domain Facets and Safety Criticality 5.1. Introductory Notions 5.1.1. Facet

- By a **domain facet** we shall understand
 - \otimes one amongst a finite set of generic ways
 - \otimes of analysing a domain:
 - \otimes a view of the domain,
 - \otimes such that the different facets cover conceptually different views,
 - \otimes and such that these views together cover the domain .

• We shall in this talk distinguish between the following facets:

∞ intrinsics,

« support technologies,

« human behaviour,

- \circledast rules & 19 regulations and
- *∞* organisation *&* management.

¹⁹We use the ampersand '&' between terms A and B to emphasize that we mean to refer to one subject, the conjoint A&B

5.1.2. Safety Criticality

• Safety critical systems

- « are those systems whose failure may result in
- *∞* the loss of life,
- « significant property damage or
- \otimes damage to the environment.²⁰

²⁰John C. Knight: Safety Critical Systems: Challenges and Directions http://www.-cs.virginia.edu/~jck/publications/knight.state.of.the.art.summary.pdf

- For each of the domain facet categories we shall look for a corresponding, domain-specific category of hazards.
 - \otimes That is, we shall view safety criticality in potentially three steps:
 - from the point of view of the domain in which a computing system is to be inserted, hence first developed,
 - ∞ from the point of view of the requirements prescribed for such a system, and
 - ∞ from the point of view of the machine (i.e., hardware + software) design of that system.
 - \otimes In this talk we shall only consider the first step.

5.2. Intrinsics

• By **domain intrinsics** we shall understand

- those phenomena and concepts of a domain which are basic to any of the other facets (listed earlier and treated, in some detail, below),
- with such domain intrinsics initially covering at least one specific, hence named, stake-holder view.
- **Example**: The introductory example focused on

 - \otimes some derived concepts (routes etc.)

• **Hazards**: The following are examples of hazards based sôlely on the intrinsics of the domain:

\otimes environmental hazards:

destruction of one or more pipeline units due to
an earth quake, an explosion, a fire or something "similar"
occurring in the immediate neighbourhood of these units;

\otimes design faults:

 ∞ the pipeline net is not acyclic;

 \otimes etcetera

- Intrinsics hazards are such which violate the well-formedness of the domain.
 - A "domain description" is presented, but it is not a well-formed domain description.
 - \otimes One could claim that
 - ∞ whichever (event) falls outside
 - the intrinsics domain description,
 - ∞ whether it violates well-formedness criteria for domain parts
 - ∞ or action, event or behaviour pre/post-conditions,
 - ∞ is a hazard.
 - \otimes In the context of system safety we shall take the position that ∞ explicitly identified hazards
 - ∞ must be described,
 - ∞ also formally.²¹

* We refer to the main example.

- * More specifically to the well-formedness of pipeline systems as expressed in wf_PLS (Slide 18, Item 2.).
- \ast We express hazards of the intrinsics of pipeline systems by named predicates over $\mathsf{PLS'}$ and not $\mathsf{PLS}.$

5.3. Support Technologies

- By domain **support technology** we shall understand
 - « technological ways and means of implementing
 - \otimes certain observed phenomena or
 - \otimes certain conceived concepts.
- The facet of support technology, as a concept,
 - « is related to actions of specific parts;
 - « that is, a part may give rise to one or more support technologies,
 - « and we say that the support technologies 'reside' in those parts.

• Examples:

- wells are, in the intrinsics facet description abstracted as atomic units but in real instances they are complicated (composite) entities of pumps, valves and pipes;
- w pumps are similarly, but perhaps not as complicated complex units;
- \otimes values likewise; and
- \otimes sinks are, in a sense, the inverse of wells \blacksquare

• Faults:

∞ a pump may fail to respond to a *stop pump* signal; and
∞ a valve may fail to respond to an *open valve* signal

• I think it is fair to say that

most papers on the design of safety critical software are on
software for the monitoring & control of support technology.

• Describing causes of errors is not simple.

²¹These tools and techniques typically include © two or more formal specification languages, for example:

| * VDM , | * Event-B, | * TLA+ and |
|------------------------------|-------------------------------|-------------------|
| * DC , | * RAISE/RSL, | * Alloy; |
| ∞ one or more theorem | n proving tools, for example: | |
| * ACL , | * Isabelle/HOL, | * PVS and |
| * Coq , | * STeP , | * Z3 ; |

 ∞ a model-checker, for example:

* **SMV** and

 ∞ and other such tools and techniques.

* SPIN/Promela;

5.4. Human Behaviour

- A proper domain description includes humans as both
 (usually atomic) parts and
 - ∞ the behaviours that we (generally) "attach" to parts.
- **Examples**: The human operators that
 - \otimes operate wells, values, pumps and sinks;
 - « check on pipeline units;
 - \otimes decide on the flow of material in pipes,
 - ⇔etcetera ■

- By domain **human behaviour** we shall understand

∞ from (i) careful, diligent and accurate,

via

- ∞ (ii) *sloppy* dispatch, and
- ∞ (iii) delinquent work,

to

 ∞ (iv) outright *criminal* pursuit.

²²— in contrast to technology

- Typically human behaviour focus on actions and behaviours that are carried out by humans.
 - « The intrinsics description of actions and behaviours
 - ∞ focus sôlely on intended, careful, diligent and accurate performance.
- **Hazards**: This leaves "all other behaviours" as hazards!
- Proper hazard analysis, however, usually
 - « explicitly identifies failed human behaviours,
 « for example, as identified deviations from
 « described actions etc.
- Hazard descriptions thus follows from "their corresponding" intrinsics descriptions

5.5. Rules & Regulations

• Rules and regulations come in pairs $(\mathcal{R}_u, \mathcal{R}_e)$.

5.5.1. **Rules**

By a domain **rule** we shall understand some text

• which prescribes how people are, or equipment is,
• "expected" (for "..." see below) to behave
• when dispatching their duty,
• respectively when performing their function.

• **Example**: There are rules for operating pumps. One is:

- \ll if there does not exist an open, embedded route $r^{\prime\prime\prime}$ such that $\langle p\rangle \widehat{\ }r^{\prime\prime\prime}$
- \otimes ends in an open sink \blacksquare

- **Hazards**: when stipulating "expected" (as above)
 - \otimes the rules more or less implicitly
 - \otimes express also the safety criticality:
 - \otimes that is, when people are, or equipment is,
 - \otimes behaving erroneously \blacksquare
- **Example**: A domain rule which states, for example,

 \otimes that a pump, p, on some well-to-sink route $r = r'^{\langle p \rangle} r''$, \otimes may be started even

- \circledast if there does not exist an open, embedded route r''' such that $\langle p\rangle \widehat{\ }r'''$
- \otimes ends in an open sink
- is a hazardous rule \blacksquare

• Modelling Rules:

- \otimes We can model a rule by giving it
 - ∞ both a syntax
 - ∞ and a semantics.
- \otimes And we can choose to model the semantics of a rule, \mathbb{R}_u ,
 - ∞ as a predicate, \mathcal{P} , over pairs of states: $\mathcal{P} : \Sigma \times \Sigma \longrightarrow \underline{\mathbf{Bool}}$.
 - ∞ That is, the meaning, \mathcal{M} , of \mathbb{R}_u is \mathcal{P} .
 - * An action or an event has changed a state σ into a state σ' .
 - * If $\mathcal{P}(\sigma, \sigma')$ is **<u>true</u>**
 - it shall mean that the rule as been obeyed.
 - * If it is $\underline{\mathbf{false}}$ it means that the rule has been violated.

5.5.2. Regulations

- By a domain **regulation** we shall understand
 - \otimes some text which "prescribe" ("...", see below)
 - \otimes the remedial actions that are to be taken
 - \otimes when it is decided
 - \otimes that a rule has not been followed
 - \otimes according to its intention .

- **Example**: There are regulations for operating pumps and valves:
 - \otimes Once it has been discovered that a rule is hazardous
 - ∞ there should be a regulation which
 - * starts an administrative procedure which ensures that the rule is replaced; and
 - ∗ starts a series of actions which somehow brings the state of the pipeline into one which poses no danger and then applies a non-hazard rule

• **Hazards**: when stipulating "prescribe"

« regulations express requirements

 \otimes to emerging hardware and software \blacksquare

• Modelling Regulations:

- \otimes We can model a regulation by giving it
 - ∞ both a syntax
 - and a semantics.
- \otimes And we can choose to model the semantics of a regulation, \mathbb{R}_e ,
 - ∞ as a state-transformer, \mathcal{S} , over pairs of states: $\mathcal{S} : \Sigma \times \Sigma \longrightarrow \Sigma$.
 - ∞ That is, the meaning, \mathcal{M} , of \mathbb{R}_e is \mathcal{S} .
 - ∞ A state-transformation $\mathcal{S}(\sigma, \sigma')$
 - ∞ for rule \mathbb{R}_u
 - ∞ results in a state σ'' where:

* if
$$\mathcal{P}(\sigma, \sigma')$$
 is true

* then
$$\sigma' = \sigma''$$
,

* else σ'' is a corrected state such that $\mathcal{P}(\sigma, \sigma'')$ is <u>true</u>.

5.5.3. Discussion

- Where do rules & regulations reside ?" That is,
 - ∞ *"Who checks that rules are obeyed ?"* and
 - ∞ "Who ensures that regulations are applied when rules fail ?"
- Are some of these checks and follow-ups relegated

 \otimes to humans (i.e., parts) or

∞ to machines (i.e., "other" parts)?

- that is, to the behaviour of part processes?
- The next section will basically answer those questions.

5.6. Organisation & Management

- To properly appreciate this section we need remind the reader of concepts introduced earlier in this talk.
 - \otimes With parts we associate
 - © mereologies, © attributes and © behaviours.
 - \otimes Support technology
 - ∞ is related to actions and
 - \otimes Humans are often modelled
 - ∞ first as parts,
 - ∞ then as their associated behaviour.

A Basis for Safety Critical Software

• these again focused on parts.

• It is out of this seeming jigsaw puzzle of

 mereologies,

 w attributes,

 w rules and
 w regulations

that we shall now form and model the concepts of

« organisation and » management.

5.6.1. Organisation

• By domain **organisation** we shall understand

- \otimes one or more partitionings of resources
 - where resources are usually representable as parts and materials and
 - ∞ where usually a resource belongs to exactly one partition;
- \otimes such that n such partitionings typically reflects
 - ∞ strategic (say partition π_s),
 - ∞ tactical (say partition π_t), respectively
 - ∞ operational (say partition π_o)
 - concerns (say for n = 3),
- \otimes and where "descending" partitions,
 - ∞ say $\pi_s, \pi_t, \pi_o,$
 - ${\scriptstyle \scriptsize \varpi}$ represents coarse, medium and fine partitions, respectively .

- **Examples**: This example only illustrates production aspects.
 - At the strategic level one may partition a pipeline system into just one component:

the entire collection of all pipeline units, π .

- « At the tactical level one may further partition the system into
 - ∞ the partition of all wells, π_{ws} ,
 - ∞ the partition of all sinks, π_{ss} , and
 - ∞ a partition of all pipeline routes, $\pi_{\ell s}$, that
 - * $\pi_{\ell s}$, is the set of all routes of π
 - * excluding wells and sinks.
- \otimes At the organisational level may further partition the system into ∞ the partitions of individual wells, π_{w_i} ($\pi_{w_i} \in \pi_{ws}$),
 - ∞ the partitions of individual sinks, π_{s_j} ($\pi_{s_i} \in \pi_{ws}$) and
 - ∞ the partitions of individual pipeline routes, π_{r_k} ($\pi_{\ell_i} \in \pi_{\ell s}$)

• A domain organisation serves

 \otimes to structure management and non-management staff levels and

- \otimes the allocation of
 - ∞ strategic,
 - $\ensuremath{\mathfrak{o}}$ tactical and
 - ∞ operational
 - concerns across all staff levels;
- \otimes and hence the "lines of command":
 - who does what, and
 - who reports to whom,
 - * administratively and
 - * functionally.

- Organisations are conceptual parts, that is,
 - \otimes partitions are concepts,
 - \otimes they are conceptual parts
 - « in addition, i.e., adjoint to physical parts.
- They serve as "place-holders" for management.

• Modelling Organisations:

- we can normally model an organisation as an attribute of some, usually composite, part.
 - Typically such a model would be in terms of the one or more partitionings of unique identifiers, π:Π, of domain parts, p:P.
 For example:
 - For example:

```
\frac{\mathbf{type}}{\mathbf{ORG} = \operatorname{Str} \times \operatorname{Tac} \times \operatorname{Ope} \times \dots}\operatorname{Str}, \operatorname{Tac}, \operatorname{Ope} = (\Pi - \underline{\mathbf{set}}) - \underline{\mathbf{set}}\underline{\mathbf{value}}\operatorname{attr}_{\mathbf{ORG}} : \mathbf{P} \to \operatorname{ORG}\underline{\mathbf{axiom}}\mathcal{P} : \operatorname{ORG} \to \dots \to \underline{\mathbf{Bool}}
```

where we leave the details of the partitionings **Str**, **Tac**, **Org**, ... and the axiom governing the individual partitionings and their relations for further analysis.

• Faults and Hazards:

- \otimes There are erroneous and there are risky organisations.
- An erroneous organisation is, for example, one in which
 one or more partitions are left isolated
 - ∞ with respect to there being no management "tow-holder".
- \otimes A $\ensuremath{\mathsf{hazardous}}$ organisation is, for example, one

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- ∞ that consists of too many partitionings,
- ∞ whereby related management becomes confused \blacksquare

A Basis for Safety Critical Software

5.6.2. Management

- By domain **management** we shall understand such people who (such decisions which)
 - \otimes determine, formulate and set standards concerning
 - strategic,tactical andoperational

decisions;

- who ensure that these decisions are passed on to (lower) levels of management, and to floor staff;
- \otimes who make sure that such orders, as they were, are indeed carried out;
- who handle undesirable deviations in the carrying out of these
 orders cum decisions;
- \otimes and who "backstops" complaints from lower management levels and from floor staff.

A Basis for Safety Critical Software

- **Example**: [Cf. examples on Slide 129].
 - \otimes At the strategic level there is the
 - ∞ overall management of the pipeline system.
 - - all sinks;
 - ∞ specific (disjoint) routes.
 - At the operational there may then be the management of
 individual wells,
 - individual sinks, and
 - ∞ individual groups of valves and pumps \blacksquare

• Modelling Management:

- ∞ Some parts are associated with strategic management.
 - ∞ They will have their unique identifiers, π : Π , belong to some partition in an **str:Str**.
- « Other parts are associated with tactical management.
 - ∞ They will have their unique identifiers, π : Π , belong to some partition in a corresponding **tac**:**Tac**.
- \otimes Yet other parts are associated with operational management.
 - ∞ They will have their unique identifiers, π : Π , belong to some partition in the corresponding **ope:Ope**.

 \otimes The "management" parts have their attributes form corresponding states (σ : Σ).

$\frac{\mathbf{type}}{\Sigma_{STR}}, \Sigma_{TAC}, \Sigma_{OPE},$

« An idealised rendition of management actions is:

value

action_{Strategic}: $\Sigma_{STR} \rightarrow \Sigma_{TAC} \rightarrow \Sigma_{OPE} \rightarrow \Sigma_{STR}$ action_{Tactical}: $\Sigma_{STR} \rightarrow \Sigma_{TAC} \rightarrow \Sigma_{OPE} \rightarrow \Sigma_{TAC}$ action_{Operational}: $\Sigma_{STR} \rightarrow \Sigma_{TAC} \rightarrow \Sigma_{OPE} \rightarrow \Sigma_{OPE}$

- $\text{ action}_{Strategic}$ expresses that strategic management considers the "global" state ($\Sigma_{STR} \times \Sigma_{TAC} \times \Sigma_{OPE}$) but potentially changes only the "strategy" state.
- \otimes action_{Tactical} expresses that tactical management considers the "global" state ($\Sigma_{STR} \times \Sigma_{TAC} \times \Sigma_{OPE}$) but potentially changes only the "tactical" state.
- \approx action_{Operational} expresses that tactical management considers the "global" state ($\Sigma_{STR} \times \Sigma_{TAC} \times \Sigma_{OPE}$) but potentially changes only the "operational" state.

- w We can normally model management as part of the behavioural model of some, usually composite part.
 - ∞ Typically such a model would be in terms communication procedures between managers, p:P, and their immediate subordinates, $\{p_1:P_1,p_2:P_2,\ldots,p_n:P_N\}$:
 - For example:

- Hazards: [Cf. faults and hazards, Slide 133.]
 - © Faults and hazards of organisations & management come about also as the result of "mis-management":
 - Strategic management updates tactical and operational management states.
 - Tactical management updates strategic and operational management states.
 - Operational management updates strategic and tactical management states.
 - ∞ That is: these states are not clearly delineated,
 - © Etcetera!

5.6.2.1 Discussion

- This section on organisation & management
 - \otimes is rather terse;
 - ∞ in fact it covers a whole, we should think,
 - \otimes novel and interesting theory
 - of business organisation & management

5.7. Discussion

- There may be other facets
 - \otimes but our point has been made:
 - ∞ that an analysis of hazards (including faults)
 - \otimes can, we think, be beneficially structured
 - \otimes by being related to reasonably distinct facets.
- A mathematical explanation of the concept of facet is needed.
 - One that helps partition the domain phenomena and concepts into disjoint descriptions.
 - \otimes We are thinking about it.

6. Conclusion

6.1. The Author's Scientific & Engineering Background

- The present author's research has since the early 1970s focused on
 - \otimes programming methodology:
 - ∞ how to develop software
 - ∞ such that it was correct
 - ∞ with respect to some specification —
 - ∞ call it requirements.
 - \otimes The emphasis was on
 - abstract software specifications and theirrefinement or transformation into code.

- \otimes Programming language semantics
 - ∞ and the stage- and step-wise development of compilers,
 - ∞ in many, up to nine stages and steps,
 - ∞ became a highlight of the 1980s.
- © The step from programming language semantics to domain descriptions followed:
 - Domain descriptions, in a sense, specified
 - ∞ the language inherent in the described domain —
 - ∞ that is: "spoken" by its actors, etc.

- Since the early 1990s I therefore additionally focused on domain descriptions.
 - w Now an additional goal of software development might be achieved:
 - \otimes securing that the software met customers' expectations.
- With the observation that
 - \otimes requirements prescriptions
 - can be systematically
 - but, of course, not automatically —
 - "derived" from domain descriptions
 - \otimes a bridge was established:
 - from domains via requirements to software.

6.2. What Have We Achieved?

- When Dr Clive Victor Boughton, on November 4, 2013, approached me on the subject of
 - « "Software Safety: New Challenges and Solutions",

 \otimes I therefore, naturally questioned:

- ∞ can one stratify the issues of safety criticality into three phases:
- ∞ searching for sources of faults and hazards in domains,
- ∞ elaborating on these while "discovering" further sources during requirements engineering, and,
- ∞ finally, during early stages of software design.
- \otimes I believe we have answered that question partially
 - ∞ with there being good hopes for further stratification.

- Yes, I would indeed claim that
 - we have contributed to the "greater" issues of safety critical systems
 - ∞ by suggesting a discipline framework for of faults "discovery" and hazards:
 - \otimes investigate the domains, the requirements and the design.

6.3. Further Work

• But, clearly, that work has only begun.

7. Acknowledgements

• I thank Dr Clive Victor Boughton of aSSCa &c.

« for having the courage to convince his colleagues to invite me,

- \otimes for having inspired me to observe that
 - faults and hazards can be "discovered"
 - purely in the context of domain descriptions,
- « for his support in answering my many questions,
- \otimes and for otherwise arranging my visit.

Thanks