From Domains to Requirements

On a Triptych of Software Development

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ABSTRACT

Domain engineering, in the sense of this paper, is offered as a means to help secure that software engineers deliver the right software - where formalisation of relevant stages and steps of software development helps secure that the software is right [10]. In this paper we shall present the essence of a software development triptych: from domains via requirements to software design. We emphasize the two first phases: domain engineering and requirements engineering. We show the pragmatic stages of the construction of domain descriptions: the facets of intrinsics, support technologies, management & organisation, rules & regulations, script (licenses and contracts) and human behaviour. And we show how to construct main facets of requirements prescriptions: domain requirements and interface requirements. In this respect we focus in particular on the domain requirements development stages of projection, instantiation, determination and extension. The paper represents a summary as well as some significant improvements over the domain-torequirements coverage of [3, Vol. 3, Parts IV-V].

Categories and Subject Descriptors

D.2 [Software Engineering]: General; D.2.1 [Software Engineering]: Requirements Engineering—derivation from domain descriptions; D.2.4 [Software Engineering]: Verification; D.2.10 [Software Engineering]: Methodologies; F.3 [Theory of Computation]: Logics and Meanings of Specification; D.3 [Software Engineering]: Specification Languages

Keywords

Domain Engineering, Requirements Engineering, Formal Specification

1. INTRODUCTION

Dogma: Before we can design software we must have a robust understanding of its requirements. And before we can prescribe requirements we must have a robust understanding of the environment, or, as we shall call it, the domain in which the software is to serve — and as it is at the time such software is first being contemplated.

In consequence we suggest that software, "ideally", be developed in three phases.

First a phase of **domain engineering.** In this phase a reasonably comprehensive description is constructed from an analysis of the domain. That description, as it evolves, is analysed with respect to inconsistencies, conflicts and completeness.

Then a phase of **requirements engineering.** This phase is strongly based, as we shall see (in Sect. 4), on an available, necessary and sufficient domain description. Guided by the domain and requirements engineers the requirements stakeholders points out which domain description parts are to be left (projected) out of the domain requirements, and of those left in what forms of instantiations, determinations and extensions are required. Similarly the requirements stakeholders, guided by the domain and requirements engineers, informs as to which domain entities, actions, events and behaviours are shared between the domain and the machine, that is, the hardware and the software being required. In this paper we shall only very briefly cover aspects of machine requirements.

And finally a phase of **software design.** We shall not cover this phase in this paper.

_ Methodology

The paper is a methodology paper — where a method is seen as a set of principles (applied by engineers, not machines) for selecting and applying (often with some tool support) techniques (and tools) for the efficient construction of some artifact - here software.

1.1 What are Domains?

By a domain we shall here understand a universe of discourse, an area of nature subject to laws of physics and

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¹Section 5 will discuss practical renditions of "idealism"!

studies by physicists, or an area of human activity subject to its interfaces with nature. There are other domains which we shall ignore. We shall focus on the human-made domains. "Large scale" examples are the financial service industry: banking, insurance, securities trading, portfolio management, etc.; health care: hospitals, clinics, patients, medical staff, etc.; transportation: road, rail/train, sea/shipping, and air/aircraft transport (vehicles, transport nets, etc.); oil and gas systems: pumps, pipes, valves, refineries, distribution, etc. "Intermediate scale" examples are automobiles: manufacturing or monitoring and control, etc.; and heating systems. The above explication was "randomised": for some domains, to wit, the financial service industry, we mentioned major functionalities, for others, to wit, health care, we mentioned major entities.

1.2 What is a Domain Description?

By a domain description we understand a description of the entities, the actions, the events and the behaviours of the domain, including its interfaces to other domains. A domain description describes the domain as it is. A domain description does not contain requirements let alone references to any software. Michael Jackson, in [18], refers to domain descriptions as indicative (stating objective fact), requirements prescriptions as optative (expressing wish or hope) and software specifications as imperative ("do it!"). A description is syntax. The meaning (semantics) of a domain description is a usually a set of domain models. We shall take domain models to be mathematical structures (theories). The form of domain descriptions that we shall advocate "come in pairs": precise, say, English text alternates with clearly related formula text.

1.3 Description Languages

Besides using as precise a subset of a national language, as here English, as possible, and in enumerated expressions and statements, we "pair" such narrative elements with corresponding enumerated clauses of a formal specification language. We shall be using the RAISE Specification Language, RSL, [12, 13, 3], in our formal texts. But any of the modeloriented approaches and languages offered by Alloy [17], Event B [1], VDM [8, 9, 11] and Z [25], should work as well. No single one of the above-mentioned formal specification languages, however, suffices. Often one has to carefully combine the above with elements of Petri Nets [21], CSP: Communicating Sequential Processes [15], MSC: Message Sequence Charts [16], Statecharts [14], and some temporal logic, for example either DC: Duration Calculus [26] or TLA+ [20]. Research into how such diverse textual and diagrammatic languages can be meaningfully and proof-theoretically combined is ongoing [2].

1.4 Contributions of This Paper

We claim that the major contributions of the Triptych approach to software engineering as presented in this paper are the following: (1) the clear identification of domain engineering, or, for some, its clear separation from requirements engineering (Sects. 3 and 4); (2) the identification and 'elaboration' of the pragmatically determined domain facets of intrinsics, support technologies, management and organisation, rules and regulations, scripts (licenses and contracts) and human behaviour whereby 'elaboration' we mean that

we provide principles and techniques for the construction of these facet description parts (Sects. 3.2–3.7); (3) the reidentification and 'elaboration' of the concept of business process re-engineering (Sect. 4.1) on the basis of the notion of business processes as first introduced in Sect. 3.1; (4) the identification and 'elaboration' of the technically determined domain requirements facets of projection, instantiation, determination, extension and fitting requirements principles and techniques — and, in particular the "discovery" that these requirements engineering stages are strongly dependent on necessary and sufficient domain descriptions (Sects. 4.2.1–4.2.5); and (5) the identification and 'elaboration' of the technically determined interface requirements facets of shared entity, shared action, shared event and shared behaviour requirements principles and techniques (Sects. 4.3.1-4.3.4). We claim that the facets of (2, 3, 4)and (5) are all novel. In Sect. 5 we shall discuss these contributions in relation to the works and contributions of other researchers and technologists.

1.5 Relation to Other Engineering Disciplines

An aeronautics engineer to be hired by Boeing to their design team for a next generation aircraft must be pretty well versed in applied mathematics and in aerodynamics. A radio communications engineer to be hired by Ericsson to their design team for a next generation mobile telephony antennas must be likewise pretty well versed in applied mathematics and in the physics of electromagnetic wave propagation in matter. And so forth. Software engineers hired for the development of software for hospitals, or for railways, know little, if anything, about health care, respectively rail transportation (scheduling, rostering, etc.). The Ericsson radio communications engineer can be expected to understand Maxwell's Equations, and to base the design of antenna characteristics on the transformation and instantiation of these equations. It is therefore quite reasonable to expect the domain-specific software engineer to understand formalisation of their domains: for railways: www.railwaydomain.org, and for pipelines pipelines.pdf, logistics logistics.pdf and for container lines container-paper.pdf - all at www.imm.dtu.dk/~db/.

1.6 Structure of Paper

Before going into some details on domain enginering (Sect. 3) and requirements engineering (Sect. 4) we shall in the next section cover the basic concepts of specifications, whether domain descriptions or requirements prescriptions. These are: entities, actions, events and behaviours. Section 5 then discuses the contributions of the Triptych approach as covered in this paper.

On Examples —

We bring 22 examples. These examples take up about 50% of the paper space. Most of these have both narrative, informal (English) text and formal texts. In principle all examples should have formal texts. But page space concerns dictated their absence.

On Formalisations

The reader, however, need not read the formalised parts of the examples! They are expressed in the RAISE [13] Specification Language (RSL [12])

2. A SPECIFICATION ONTOLOGY

In order to describe domains we postulate the following related specification components: entities, actions, events and behaviours.

2.1 Entities

By an entity we shall understand a phenomenon we can point to in the domain or a concept formed from such phenomena

Example 1. *Entities:* The example is that of aspects of a transportation net. You may think of such a net as being either a road net, a rail net, a shipping net or an air traffic net. Hubs are then street intersections, train stations, harbours, respectively airports. Links are then street segments between immediately adjacent intersections, rail tracks between train stations, sea lanes between harbours, respectively air lanes between airports.

- 1. There are hubs and links.
- There are nets, and a net consists of a set of two or more hubs and one or more links.

```
type

1 H, L,

2 N = H-set \times L-set
axiom

2 \forall (hs,ls):N • card hs\geq2 \wedge card ks\geq1
```

- 3. There are hub and link identifiers
- 4. Each hub (and each link) has an own, unique hub (respectively link) identifier (which can be observed (ω) from the hub [respectively link]).

```
type
3 HI, LI
value
4a \omegaHI: H \rightarrow HI, \omegaLI: L \rightarrow LI
axiom
4b \forall h,h':H, I,I':L \bullet h\neqh' \Rightarrow
\omegaHI(h)\neq\omegaHI(h') \wedge
|\neq|'\Rightarrow\omegaLI(I)\neq\omegaLI(I')
```

In order to model the physical (i.e., domain) fact that links are delimited by two hubs and that one or more links emanate from and are, at the same time incident upon a hub we express the following:

- From any link of a net one can observe the two hubs to which the link is connected.
 - (a) We take this 'observing' to mean the following: From any link of a net one can observe the two distinct identifiers of these hubs.
- 6. From any hub of a net one can observe the one or more links to which are connected to the hub.
 - (a) Again: by observing their distinct link identifiers.
- 7. Extending Item 5: the observed hub identifiers must be identifiers of hubs of the net to which the link belongs.
- 8. Extending Item 6: the observed link identifiers must be identifiers of links of the net to which the hub belongs

We used, above, the concept of 'identifiers of hubs' and 'identifiers of links' of nets. We define, below, functions (iohs, iols) which calculate these sets.

```
value Formalisation _____
```

```
5a \omegaHIs: L \rightarrow HI-set,
      6a \omega Lls: H \rightarrow Ll\text{-set}.
axiom
      5b \forall I:L • card \omegaHIs(I)=2 \wedge
      6b \forall h:H • card \omegaLls(h)\geq1 \wedge
      \forall (hs,ls):N •
                   \forall h:H • h \in hs \Rightarrow
      5a)
                        \forall li:Ll • li \in \omegaLls(h) \Rightarrow
                              \exists l': L \bullet l' \in ls \land li = \omega Ll(l') \land
                                    \omega HI(h) \in \omega HIs(I') \wedge
                    \forall \text{ I:L } \bullet \text{ I} \in \text{ Is } \Rightarrow
      6a)
                        \exists h',h''':H \bullet \{h',h''\}\subseteq hs \land
                             \omega HIs(I) = \{\omega HI(h'), \omega HI(h'')\}
      7 \forall h:H • h \in hs \Rightarrow \omega Lls(h) \subseteq iols(ls)
      8 \forall I:L • I \in Is \Rightarrow \omegaHIs(h) \subseteq iohs(hs)
value
      iohs: H-set \rightarrow HI-set, iols: L-set \rightarrow LI-set
      iohs(hs) \equiv \{\omega HI(h)|h:H \bullet h \in hs\}
      iols(ls) \equiv \{\omega LI(l)|l:L \bullet l \in ls\}
```

In the above extensive example we have focused on just five entities: nets, hubs, links and their identifiers. The nets, hubs and links can be seen as separable phenomena. The hub and link identifiers are conceptual models of the fact that hubs and links are connected — so the identifiers are abstract models of 'connection', or, as we shall later discuss it, the mereology of nets, that is, of how nets are composed. These identifiers are attributes of entities. Links and hubs have been modelled to possess link and hub identifiers. A link's "own" link identifier enables us to refer to the link, A link's two hub identifiers enables us to refer to the connected hubs. Similarly for the hub and link identifiers of hubs.

- 9. A hub, h_i , state, $h\sigma$, is a set of hub traversals.
- 10. A hub traversal is a triple of link, hub and link identifiers $(l_{i_{in}},h_{i_{i}},l_{i_{out}})$ such that $l_{i_{in}}$ and $l_{i_{out}}$ can be observed from hub h_{i} and such that $h_{i_{i}}$ is the identifier of hub h_{i} (wf_H Σ).
- 11. A hub state space is a set of hub states such that all hub states concern the same hub ($wf_-H\Omega$).

```
type
       9 HT = (LI \times HI \times LI)
       10 H\Sigma = HT-set
       11 H\Omega = H\Sigma-set
value
       10 \omega H\Sigma: H \rightarrow H\Sigma
       11 \omega H\Omega: H \rightarrow H\Omega
axiom
       \forall n:N,h:H•h \in \omegaHs(n) \Rightarrow
              wf_H\Sigma(\omega H\Sigma(h)) \wedge wf_H\Omega(h,\omega H\Omega(h))
       \mathsf{wf\_H}\Sigma \colon \mathsf{H}\Sigma \to \mathbf{Bool}, \, \mathsf{wf\_H}\Omega \colon \mathsf{H} \times \mathsf{H}\Omega \to \mathbf{Bool}
       wf_H\Sigma(h\sigma) \equiv
              \forall (li,hi,li'),(\underline{\phantom{a}},hi',\underline{\phantom{a}}):\mathsf{HT}\bullet(li,hi,li')\in \mathsf{h}\sigma \Rightarrow
                            \{li,li'\}\subseteq\omega Lls(h)\wedge hi=\omega Hl(h)\wedge hi'=hi
       wf_H\Omega(h,h\omega) \equiv
              \forall \ \mathsf{h}\sigma : \mathsf{H}\Sigma \bullet \mathsf{h}\sigma \in \mathsf{h}\omega \Rightarrow \mathsf{wf}_{\mathsf{L}}\mathsf{H}\Sigma (\mathsf{h}\sigma) \wedge \mathsf{h}\sigma \neq \{\}
                     \Rightarrow let (li,hi,li'):HT•(li,hi,li')\in h\sigma in hi=\omegaHI(h) end
```

■ End of Example

2.2 Actions

A set of entities are form a state. It is the domain engineer which decides on such states. A function, f, application, one which applies to zero, one or more arguments and a state results in a state changes, is an action.

Example 2. Actions:

- 12. Our example action is that of setting the state of hub.
- 13. The setting applies to a hub
- 14. and a hub state in the hub state space
- 13. and vields a "new" hub.
- 15. The before and after hub identifier remains the same.
- 16. The before and after hub state space remains the same.
- 17. The result hub is in the hub state space subject to some probability distribution.

```
value

p:Real, axiom 0<p≤1, typically p≈ 1 − 10<sup>-7</sup>

\overline{p}:Real, axiom \overline{p}=1−p

12 set_HΣ: H × HΣ → H

13 set_HΣ(h,hσ) as h'

14 pre hσ ∈ ωHΩ(h)

15 post ωHI(h)=ωHI(h')∧

16 ωHΩ(h)=ωHΩ(h')∧

15 ωHΣ(h')=

15 ([{hσ'|hσ':HΣ•hσ'∈ ωΩ(h)\{hσ}})_{\overline{p}}]_phσ
```

The non-deterministic internal choice operator expression $s_{\overline{p}} \lceil p s'$ with probability p has value s' and with probability \overline{p} has value s. The prefix internal choice operator expression $\lceil \{h\sigma_i, h\sigma_j, \dots, h\sigma_k\}$ h non-deterministically as one of the values in the set $\{h\sigma_i, h\sigma_j, \dots, h\sigma_k\}$, that is, is the same as $h\sigma_i \lceil h\sigma_j \rceil \dots \lceil h\sigma_k \rceil$. End of Example

2.3 Events

Any state change is an event. A situation in which a (specific) state change was expected but none (or another) occurred is an event. Some events are more "interesting" than other events. Not all state changes are caused by actions of the domain.

Example $\ 3.$ Events: We keep to our ongoing example of transportation nets.

- 18. A hub is in some state, $h\sigma$.
- 19. An action directs it to change to state $h\sigma'$ where $h\sigma' \neq h\sigma$.
- 20. But after that action the hub remains either in state $h\sigma$ or is possibly in a third state, $h\sigma''$ where $h\sigma'\neq h\sigma''$.
- 21. Thus an "interesting event" has occurred!

```
Formalisation
\exists \ n:N,h:H,h\sigma,h\sigma':H\Sigma^{\bullet}h\in\omega Hs(n)\wedge
19,20 \ \{h\sigma,h\sigma'\}\subseteq\omega H\Omega(h)\wedge \mathbf{card}\{h\sigma,h\sigma'\}=2 \wedge
18 \ \omega H\Sigma(h)=h\sigma ;
19 \ let \ h'= \mathsf{set\_H\Sigma}(h,h\sigma') \ in
20 \ \omega H\Sigma(h')\in\omega H\Sigma(h')\backslash\{h\sigma'\}\Rightarrow
21 \ ''interesting \ event'' \ end
```

It only makes sense to change hub states if there are more than just one single such state. $\hfill\blacksquare$ End of Example

2.4 Behaviours

A behaviour is a set of zero, one or more sequences of actions, including events.

Example $4.\ Behaviours:$ We keep to our ongoing example of transportation net changes.

- 22. Let h be a hub of a net n.
- 23. Let $h\sigma$ and $h\sigma'$ be two distinct states of h.

- 24. Let ti:TI be some time interval, say 3sec.
- 25. Let h start in an initial state $h\sigma$.
- 26. Now subject hub h to an ongoing sequence of, say, 100 state changes
 - (a) from $h\sigma$ to $h\sigma'$ and
 - (b) then, after a wait of ti seconds,
 - (c) to $h\sigma$.

```
Formalisation
type
     ΤI
value
     ti:TI
     n:Nat
     26 blinking: H \times H\Sigma \times H\Sigma \rightarrow H
     26 blinking(h,h\sigma,h\sigma',m) in
             let h' = \text{set\_H}\Sigma(h, h\sigma) in
     25
     26c
              wait ti;
              let h'' = set_H\Sigma(h',h\sigma') in
     26a
              wait ti;
              \mathbf{if}\;m{=}1\;\mathbf{then}\;h''
     26
              else blinking(h,h\sigma,h\sigma',m-1) end end end
     26
              \mathbf{pre}\ \{\mathsf{h}\sigma,\mathsf{h}\sigma'\}{\subseteq}\omega\mathsf{H}\Omega(\mathsf{h}){\wedge}\mathsf{h}\sigma{\neq}\mathsf{h}\sigma'
     23
                    \land initial m=100
```

■ End of Example

3. DOMAIN ENGINEERING

We focus on the facet components of a domain description and leave it to other publications, for ex. [3, Vol. 3, Part IV, Chaps. 8–10], to cover such aspects of domain engineering as stake-holder identification and liaison, domain acquisition and analysis, terminologisation, verification, testing, model-checking, validation and domain theory. By understanding, first, the facet components the domain engineer is in a better position to effectively establish the regime of stakeholders, pursue acquisition and analysis, and construct a necessary and sufficient terminology. The domain description components each cover their domain facet. We outline six such facets: intrinsics, support technology, rules and regulations, scripts (licenses and contracts), management and organisation, and human behaviour. But first we cover a notion of business processes.

3.1 Business Processes

By a business process we understand a set of one or more, possibly interacting behaviours which fulfill a business objective. We advocate that domain engineers, typically together with domain stake-holder groups, rough-sketch their individual business processes.

Example 5. Some Transport Net Business Processes: With respect to one and the same underlying road net we suggest some business-processes and invite the reader to rough-sketch these.

- 27. Private citizen automobile transports: Private citizens use the road net for pleasure and for business, for sightseeing and to get to and from work.
- 28. Public bus (&c.) transport: Province and city councils contract bus (&c.) companies to provide regular person transports according to timetables and at cost or free of cost.

- 29. Road maintenance and repair: Province and city councils contract service companies to monitor road (link and hub) surface quality, to maintain set standards of surface quality, and to "emergency" re-establish sudden occurrences of low standards.
- 30. *Toll road traffic:* State government and province councils contract companies to run toll road nets with toll booth plazas.
- 31. Net revision: road (&c.) building: State government and province and city councils contract road building contractors to extend (or shrink) road nets, i.e., inserting new or removing old hubs and links.

The detailed narrative and formalisation of the above rough-sketched business process synopses now becomes part of the domain description as partially exemplified in the previous and the next many examples.

■ End of Example

Rough-sketching such business processes helps bootstrap the process of domain acquisition. We shall return to the notion of business processes in Sect. 4.1 where we introduce the concept of business process re-engineering.

3.2 Intrinsics

By intrinsics we shall understand the very basics, that without which none of the other facets can be described, i.e., that which is common to two or more of these other facets.

Example 6. *Intrinsics:* Most of the descriptions of Sect. 2 model intrinsics. We add a little more.

- 32. A link traversal is a triple of a (from) hub identifier, an along link identifier, and a (towards) hub identifier
- 33. such that these identifiers make sense in any given net.
- 34. A link state is a set of link traversals.
- 35. And a link state space is a set of link states.

```
\label{eq:type} \begin{tabular}{ll} type & & & & & \\ 32 \ LT = HI \times LI \times HI \\ & & & & & \\ 34 \ L\Sigma' = LT\text{-set} \\ & & & & \\ 34 \ L\Sigma = \{||\sigma:L\Sigma'\bullet \mathsf{wf}\_L\Sigma(|\sigma)|\} \\ & & & & \\ 35 \ L\Omega' = L\Sigma\text{-set} \\ & & & \\ 35 \ L\Omega = \{||\omega:L\Omega'\bullet \mathsf{wf}\_L\Omega(|\omega)|\} \\ & & & \\ value & & & \\ 33 \ \mathsf{wf}\_LT: LT \to \mathsf{N} \to \mathbf{Bool} \\ & & & \\ 33 \ \mathsf{wf}\_LT(\mathsf{hi},\mathsf{li},\mathsf{hi}')(\mathsf{n}) \equiv \\ & & & \\ 33 \ \mathsf{wf}\_LT(\mathsf{hi},\mathsf{li},\mathsf{hi}')(\mathsf{n}) \equiv \\ & & & \\ 33 \ \mathsf{wf}\_LT(\mathsf{hi},\mathsf{li},\mathsf{hi}')(\mathsf{n}) \equiv \\ & & & \\ 33 \ \mathsf{wf}\_LT(\mathsf{hi},\mathsf{li},\mathsf{hi}')(\mathsf{n}) \equiv \\ & & & \\ 33 \ \mathsf{wf}\_LT(\mathsf{hi},\mathsf{li},\mathsf{hi}')(\mathsf{n}) \equiv \\ & & & \\ 33 \ \mathsf{wf}\_LT(\mathsf{hi},\mathsf{li},\mathsf{hi}')(\mathsf{n}) \equiv \\ & & & \\ 33 \ \mathsf{wf}\_LT(\mathsf{hi},\mathsf{li},\mathsf{hi}')(\mathsf{n}) \equiv \\ & & \\ 33 \ \mathsf{wf}\_LT(\mathsf{hi},\mathsf{li},\mathsf{li},\mathsf{li}')(\mathsf{n}) \equiv \\ & & \\ 33 \ \mathsf{wf}\_LT(\mathsf{hi},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li}')(\mathsf{n}) \equiv \\ & & \\ 33 \ \mathsf{wf}\_LT(\mathsf{hi},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li})(\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li}')(\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},\mathsf{li},
```

The wf_L Σ and wf_L Ω can be defined like the corresponding functions for hub states and hub state spaces.

■ End of Example

The intrinsics facet has an own theory and anomy develop.

The *intrinsics* facet has an own theory and anown development methodology.

3.3 Support Technologies

By support technologies we shall understand the ways and means by which humans and/or technologies support the representation of entities and the carrying out of actions.

Example 7. Support Technologies: Some road intersections (i.e., hubs) are controlled by semaphores alternately shining red, yellow, green in carefully interleaved sequences in each of the indirections from links incident upon the hubs. Usually these signalings are initiated as a result of road traffic sensors placed below the surface of these links. We shall model just the signaling.

- 36. There are three colours: red, yellow and green.
- Each hub traversal is extended with a colour and so is the hub state.

- 38. There is a notion of time interval.
- 39. Signaling is now a sequence, $\langle (h\sigma',t\delta'), (h\sigma'',\ t\delta''), \ldots, (h\sigma'\cdots',\ t\delta'\cdots') \rangle$ such that the first hub state $h\sigma'$ is to be set first and followed by a time delay $t\delta'$ whereupon the next state is set, etc.
- 40. A semaphore is now abstracted by the signalings that are prescribed for any change from a hub state $h\sigma$ to a hub state $h\sigma'$.
- 41. We assume an idealised version of the chg_H Σ action.

```
Formalisation _
type
    36 Colour == red | yellow | green
    37 X = LI \times HI \times LI \times Colour [crossings of a hub]
    37 H\Sigma = X-set [hub states]
    38 TI [time interval]
    39 Signalling = (H\Sigma \times TI)^*
    40 Semaphore = (H\Sigma \times H\Sigma) \overrightarrow{m} Signalling
value
    37 \omega H\Sigma: H \rightarrow H\Sigma
    40 \omegaSemaphore: H \rightarrow Sema,
    41 chg_H\Sigma: H \times H\Sigma \rightarrow H
    41 chg_H\Sigma(h,h\sigma) as h'
           \mathbf{pre}\ \mathsf{h}\sigma \in \omega \mathsf{H}\Omega(\mathsf{h})\ \mathbf{post}\ \omega \mathsf{H}\Sigma(\mathsf{h}') {=} \mathsf{h}\sigma
    41
    39 chg_H\Sigma_Seq: H \times H\Sigma \rightarrow H
    39 chg_H\Sigma_Seq(h,h\sigma) \equiv
            let sigseq = (\omega Semaphore(h))(\omega \Sigma(h),h\sigma) in
            sig_seq(h)(sigseq) end
          \mathsf{sig\_seq} \colon \mathsf{H} \to \mathsf{Signalling} \to \mathsf{H}
    39
          sig\_seq(h)(sigseq) \equiv
            if \mathsf{sigseq} {=} \langle \rangle then \mathsf{h} else
    39
             let (h\sigma,t\delta) = hd sigseq in let h' = chg\_H\Sigma(h,h\sigma);
    39
             wait t\delta:
            sig_seq(h')(tl sigseq) end end end
```

■ End of Example

The support technologies facet has an own theory and anown development methodology.

3.4 Rules and Regulations

By a rule we shall understand a text which describe how the domain is (i.e., people and technology are) expected to behave. The meaning of a rule is a predicate over "before/after" states of actions (simple, one step behaviours): if the predicate holds then the rule has been obeyed. By a regulation we shall understand a text which describes actions to be performed should its corresponding rule fail to hold. The meaning of a regulation is therefore a state-to-state transition, one that brings the domain into a rule-holding "after" state.

Example $8.\ Rules:$ We give two examples related to railway systems where train stations are the hubs and the rail tracks between train stations are the links.

- 42. Trains arriving at or leaving train stations:
 - (a) (In China:) No two trains
 - (b) must arrive at or leave a train station
 - (c) in any two minute time interval.
- 43. Trains travelling "down" a railway track. We must introduce a notion of links being a sequence of adjacent sectors.
 - (a) Trains must travel in the same direction;
 - (b) and there must be at least one "free-from-trains" sector
 - (c) between any two such trains.

We omit showing somewhat "lengthy" formalisations.

■ End of Example

We omit exemplification of regulations.

The rules and regulations facet has an own theory and anown development methodology.

3.5 Scripts, Licenses and Contracts

By a script we understand a set of pairs of rules and regulations.

Example 9. Timetable Scripts:

- 44. Time is considered discrete. Bus lines and bus rides have unique names (across any set of time tables).
- 45. A TimeTable associates Bus Line Identifiers to sets of Journies.
- Journies are designated by a pair of a BusRoute and a set of BusRides.
- 47. A BusRoute is a triple of the Bus Stop of origin, a list of zero, one or more intermediate Bus Stops and a destination Bus Stop.
- 48. A set of BusRides associates, to each of a number of Bus Identifiers a Bus Schedule.
- 49. A Bus Schedule is a triple of the initial departure Time, a list of zero, one or more intermediate bus stop Times and a destination arrival Time.
- 50. A Bus Stop (i.e., its position) is a Fraction of the distance along a link (identified by a Link Identifier) from an identified hub to an identified hub.
- 51. A Fraction is a \mathbf{Real} properly between 0 and 1.
- 52. The Journies must be well_formed in the context of some net.

 Formalisation _____

type 44. T, BLId, BId 45. TT = BLId m Journies 46. Journies' = BusRoute × BusRides 47. BusRoute = BusStop × BusStop* × BusStop 48. BusRides = BId m BusSched 49. BusSched = T × T* × T 50. BusStop == mkBS(s_fhi:HI,s_ol:LI,s_f:Frac,s_thi:HI) 51. Frac = {|r:Real•0<r<1|} 52. Journies = {|j:Journies'•∃ n:N • wf_Journies(j)(n)|}

- 53. A set of journies is well-formed
- 54. if the bus stops are all different,
- 55. if a bus line is embedded in some line of the net, and

commensurable_bus_trips((bs1,bsl,bsn),js)(hs,ls)

56. if all defined bus trips of a bus line are commensurable.

value 53. wf_Journies: Journies → N → Bool 53. wf_Journies((bs1,bs1,bsn),js)(hs,ls) ≡ 54. diff_bus_stops(bs1,bsl,bsn) \land 55. is_net_embedded_bus_line($(bs1)^b$)(hs,ls) \land

■ End of Example

Timetables are used in the next example.

56

By a license (a contract) language we understand a pair of languages of licenses and of the set of actions allowed by the license – such that non-allowable actions incur moral obligations whereas, for contracts, they incur legal responsibilities.

Example 10. *Contracts:* An example contract can be 'schematiced':

```
cid: contractor cor contracts sub-contractor cee
  to perform operations
    {"conduct","cancel","insert","subcontract"}
    with respect to timetable tt.
```

We assume a context (a global state) in which all contract actions (including contracting) takes place and in which the implicit net is defined.

- 57. Contracts, contractors and sub-contractors have unique identifiers Cld, CNm, CNm.
- 58. A contract has a unique identification, names the contractor and the sub-contractor (and we assume the contractor and sub-contractor names to be distinct). A contract also specifies a contract body.
- 59. A contract body stipulates a timetable and the set of operations that are mandated or allowed by the contractor.
- 60. An Operation is either a "conduct" (i.e., conduct a bus ride), a bus ride "cancel"lation, a bus ride "insert", or a [sub]-"contract"ing operation.

```
type
57. Cld, CNm
58. Contract = Cld × CNm × CNm × Body
59. Body = Op-set × TT
60. Op == "conduct" | "cancel" | "insert" | "contract"
```

Concrete examples of actions can be schematised:

- (a) cid: conduct bus ride (blid,bid) to start at time t
- (b) cid: cancel bus ride (blid,bid) at time t
- (c) cid: insert bus ride like (blid,bid) at time t

The schematised license shown earlier is almost like an action; here is the action form:

```
(d) cid: contractor cnm' is granted a contract cid' to perform operations
{"conduct","cancel","insert",sublicense"}
with respect to timetable tt'.
```

All actions are being performed by a sub-contractor in a context which defines that sub-contractor cnm, the relevant net, say n, the base contract, referred here to by cid (from which this is a sublicense), and a timetable tt of which tt' is a subset. contract name cnm' is new and is to be unique. The subcontracting action can (thus) be simply transformed into a contract as shown on PageS-lide .

```
\label{eq:type} \textbf{Lype} \\ \textbf{Action} = \textbf{CNm} \times \textbf{Cld} \times (\textbf{SubCon} \mid \textbf{SmpAct}) \times \textbf{Time} \\ \textbf{SmpAct} = \textbf{Conduct} \mid \textbf{Cancel} \mid \textbf{Insert} \\ \textbf{Conduct} = \mu \textbf{Conduct}(\textbf{s\_blid:BLld,s\_bid:Bld}) \\ \textbf{Cancel} == \mu \textbf{Cancel}(\textbf{s\_blid:BLld,s\_bid:Bld}) \\ \textbf{Insert} = \mu \textbf{Insert}(\textbf{s\_blid:BLld,s\_bid:Bld}) \\ \textbf{SubCon} == \mu \textbf{SubCon}(\textbf{s\_cid:Cld,s\_cnm:CNm,s\_body:body}) \\ \textbf{where body} = (\textbf{s\_ops:Op-set,s\_tt:TT}) \\ \end{cases}
```

We omit formalising the semantics of these syntaxes. Our formalisation would be expressed (in CSP [15]) with each bus, each licensee (and licensor), time and the road net bus traffic being processes, etc.

■ End of Example

The scripts, licenses and contracts facet has an own theory and anown development methodology.

3.6 Management and Organisation

By management we shall understand the set of behaviours which perform strategic, tactical and operational actions. By organisation we shall understand the decomposition of these behaviours into, for example, clearly separate strategic, tactical and operational "areas", possibly further decomposed by geographical and/or "subject matter" concerns. To explain differences between strategic, tactical and operational issues we introduce a notion of strategic, tactical and operational funds, $\mathbb{F}_{S,\mathcal{T},\mathcal{O}}$, and other resources, \mathbb{R} , a notion of contexts, \mathbb{C} , and a notion of states, \mathbb{S} . Contexts bind resources to bindings from locations to disjoint time intervals (allocation and scheduling), states bind resource identifiers to resource values. Simplified types of the strategic, tactical and operational actions are now:

```
type \begin{array}{l} \mathbb{R}, \mathbb{RID}, \mathbb{RVAL}, \mathbb{F}_{\mathcal{S}}, \mathbb{F}_{\mathcal{T}}, \mathbb{F}_{\mathcal{O}} \\ \mathbb{C} = \mathbb{R} \overset{}{\text{re}} \left( (\mathbb{T} \times \mathbb{T}) \overset{}{\text{re}} \mathbb{L} \right) \\ \mathbb{S} = \mathbb{RID} \overset{}{\text{re}} \mathbb{RVAL} \\ \textbf{value} \\ \\ \omega \mathbb{RID} : \mathbb{R} \to \mathbb{RID} \\ \omega \mathbb{RVAL} : \mathbb{R} \to \mathbb{RVAL} \\ \text{Executive\_functions: } \mathbb{C} \times \mathbb{S} \times \mathbb{F}_{\mathcal{S}, \mathcal{T}, \mathcal{O}} \to \mathbb{F}_{\mathcal{S}, \mathcal{T}, \mathcal{O}} \\ \text{Strategic\_functions: } \mathbb{C} \times \mathbb{S} \times \mathbb{F}_{\mathcal{S}} \to \mathbb{F}_{\mathcal{S}} \times \mathbb{R} \times \mathbb{C} \times \mathbb{S} \\ \text{Tactic\_functions: } \mathbb{R} \times \mathbb{C} \times \mathbb{S} \times \mathbb{F}_{\mathcal{T}} \to \mathbb{C} \times \mathbb{F}_{\mathcal{T}} \\ \text{Operational\_functions: } \mathbb{C} \times \mathbb{S} \times \mathbb{F}_{\mathcal{O}} \to \mathbb{S} \times \mathbb{F}_{\mathcal{O}} \\ \end{array}
```

where we have omitted arguments pertinent to specific functions. Executive functions redistribute financial assets.

Example 11. Management: We relate to Example 10.

- The conduct, cancel and insert bus ride actions are operational functions.
- 62. The actual subcontract actions are tactical functions;
- 63. but the decision to carry out such a tactical function may very well be a strategic function as would be the acquisition or disposal of busses.
- 64. Forming new timetables, in consort with the contractor, is a strategic function.

We omit formalisations.

■ End of Example

The management and organisation facet has an own theory and anown development methodology.

3.7 Human Behaviour

By human behaviour we shall understand those aspects of the behaviour of domain stake-holders which have a direct bearing on the "functioning" of the domain, in a spectrum from diligent via sloppy to delinquent and outright criminal neglect in the observance of maintaining entities, carrying our actions and responding to events.

Example 12. *Human Behaviour:* We relate to Examples 10–11

- 65. no failures to conduct a bus ride must be classified as diligent;
- 66. rare failures to conduct a bus ride must be classified as sloppy if no technical reasons were the cause;
- 67. occasional failures · · · as delinquent;
- 68. repeated patterns of failures · · · as criminal.

We omit formalisations.

■ End of Example

The human behaviour facet has an own theory and anown development methodology.

3.8 Discussion

We have ever so briefly outlined the six concepts of domain facets and we have exemplified each of these. Real-scale domain descriptions are, of course, much larger than what we can show. Typically, say for the domain of logistics, a description is 30 pages; for "small" parts of railway systems we easily get up to 100-200 pages of the kind shown here. The reader should now have gotten a reasonably clear idea as to what constitutes a domain description. As mentioned, in the introduction to Sect. 3, we shall not cover post-modelling activities such a validation and domain theory formation. The latter is usually part of the verification (theorem proving, model checking and formal testing) of the formal domain description. Final validation of a domain description is with respect to the narrative part of the narrative/formalisation pairs of descriptions. The reader should also be able to form a technical opinion about what can be formalised, and that not all can be formalised within the framework of a single formal specification language, cf. Sect. 1.3.

4. REQUIREMENTS ENGINEERING

Whereas a domain description presents a domain as it is, a requirements prescription presents a domain as it would be if some required machine was implemented (from these requirements). The machine is the hardware plus software to be designed from the requirements. That is, the *machine* is what the requirements are about. We distinguish between three kinds of requirements: (Sect. 4.2) the domain requirements are those requirements which can be expressed solely using terms of the domain; (Sect. 4.4) the machine requirements are those requirements which can be expressed solely using terms of the machine and (Sect. 4.3) the interface requirement are those requirements which must use terms from both the domain and the machine in order to be expressed.

4.1 Business Process Re-engineering

In Sect. 3.1 we very briefly covered a notion of business processes. These were the business processes of the domain before installation of possible computing systems. The potential of installing computing systems invariably requires revision of established business processes. Business process re-engineering (BPR) is a development of new business processes — whether or not complemented by computing and communication. BPR, such as we advocate it, proceeds on the basis of an existing domain description and outlines needed changes (additions, deletions, modifications) to entities, actions, events and behaviours following the six domain facets outlined in Sects. 3.2–3.7.

Example 13. Rough-sketching a Re-engineered Road Net: Our sketch centers around a toll road net with toll booth plazas. The BPR focuses first on entities, actions, events and behaviours (Sect. 2), then on the six domain facets (Sects. 3.2-3.7). Re-engineered Entities: Instead of the very general nets of links and hubs we shall focus on the specific road nets of toll roads. And we shall, in particular focus on a linear sequence of toll road intersections (i.e., hubs) connected by pairs of one-way (opposite direction) toll roads (i.e., links). Each toll road intersection is connected by a two way road to a toll plaza. Each toll plaza contains a pair of sets of entry and exit toll booths. (Example 15 brings more details.) Re-engineered Actions: Cars enter and leave the toll road net through one of the toll plazas. Upon entering, car drivers receive, from the entry booth, a plastic/paper/electronic ticket which they place in a special holder in the front window. Cars arriving at intermediate toll road intersections choose, on their own, to turn either "up" the toll road or "down" the toll road — with that choice being

registered by the electronic ticket. Cars arriving at a toll road intersection may choose to "circle" around that intersection one or more times — with that choice being registered by the electronic ticket. Upon leaving, car drivers "return" their electronic ticket to the exit booth and pay the amount "asked" for. Re-engineered Events: A car entering the toll road net at a toll both plaza entry booth constitutes an event. A car leaving the toll road net at a toll both plaza entry booth constitutes an event. A car entering a toll road hub constitutes an event. A car entering a toll road link constitutes an event. Re-engineered Behaviours: The journey of a car, from entering the toll road net at a toll booth plaza, via repeated visits to toll road intersections interleaved with repeated visits to toll road links to leaving the toll road net at a toll booth plaza, constitutes a behaviour — with receipt of tickets, return of tickets and payment of fees being part of these behaviours. Notice that a toll road visitor is allowed to cruise "up" and "down" the linear toll road net while (probably) paying for that pleasure (through the recordings of "repeated" hub and link entries). Re-engineered Intrinsics: No need. Re-engineered Support Technologies: There is a definite need for domain describing the toll plaza entry and exit booths. There was no prior "technology" - other than driving from a link onto a hub, now a plaza! Re-engineered Manage-- other than driving ment and Organisation: There is a definite need for domain describing the management and possibly distributed organisation of toll booth plazas. Re-engineered Rules and Regulations: Rules for entering and leaving toll booth entry and exit booths must be described as must related regulations. Rules and regulations for driving around the toll road net must be likewise be described. Re-engineered Scripts: No need. Re-engineered Human Behaviour: Humans, in this case car drivers, may not change their behaviour in the spectrum from diligent and accurate via sloppy and delinquent to outright traffic-law breaking so we see no need for any "re-engineering". ■ End of Example

The business processes facet has an own theory and anown development methodology.

4.2 Domain Requirements

For the phase of domain requirements the requirements stakeholders "sit together" with the domain cum requirements engineers and read the domain description, line-by-line, in order to "derive" the domain requirements. They do so in five rounds (in which the BPR rough sketch is both regularly referred to and possibly, i.e., most likely regularly updated). Domain requirements are "derived" from the domain description as covered in Sect. 4.2.1–4.2.5.

4.2.1 Projection

By domain projection we understand an operation that applies to a domain description and yields a domain requirements prescription. The latter represents a projection of the former in which only those parts of the domain are present that shall be of interest in the ongoing requirements development

Example 14. *Projection:* Our requirements is for a simple toll road: a linear sequence of links and hubs such as outlined in Example 13:

```
\begin{tabular}{ll} type & & & \\ N, L, H, LI, HI \\ value & & & \\ \omega Hs: \ N \rightarrow H\text{-set}, \ \omega Ls: \ N \rightarrow L\text{-set} \\ \omega HI: \ H \rightarrow HI, \qquad \omega LI: \ L \rightarrow LI \\ \omega HIs: \ L \rightarrow HI\text{-set}, \ \omega LIs: \ H \rightarrow LI\text{-set} \\ axiom & \\ See \ Items \ 5-8 \\ \end{tabular}
```

■ End of Example

The projection facet has an own theory and anown development methodology.

4.2.2 Instantiation

By domain instantiation we understand an operation that applies to a (projected) domain description, i.e., a requirements prescription, and yields a domain requirements prescription, where the latter has been made more specific, usually by constraining a domain description

Example 15. Instantiation: Here the toll road net topology is introduced.

```
type

H, L, P = H

N' = (H \times L) \times H \times ((L \times L) \times H \times (H \times L))^*

N'' = \{|n:N'' \bullet wf(n)|\}

value

wf_-N'': N' \to \mathbf{Bool}

wf_-N''((h,l),h',llhpl) \equiv ... 6 \text{ lines } ... !

\alpha N: N'' \to N

\alpha N((h,l),h',llhpl) \equiv ... 2 \text{ lines } ... !
```

wf_N" secures linearity; αN allows abstraction from more concrete N" to more abstract N.

• End of Example The *instantiation* facet has an own theory and anown development methodology.

4.2.3 Determination

By domain determination we understand an operation that applies to a (projected and possibly instantiated) domain description, i.e., a requirements prescription, and yields a domain requirements prescription, where (attributes of) entities, actions, events and behaviours have been made less indeterminate.

Example 16. *Determination:* All links and all non-end hubs are open in both directions; we leave end-hub states undefined — but see below, under 'Extension'.

```
type \\ L\Sigma = (HI\times HI)\text{-set}, L\Omega = L\Sigma\text{-set} \\ H\Sigma = (LI\times LI)\text{-set}, H\Omega = H\Sigma\text{-set} \\ N' = (H\times L)\times H\times ((L\times L)\times H\times (H\times L))^* \\ value \\ \omega L\Sigma: L\to L\Sigma, \omega L\Omega: L\to L\Omega \\ \omega H\Sigma: H\to H\Sigma, \omega H\Omega: H\to H\Omega \\ axiom \\ \forall ((h,l),h',llhhl:\langle(l',l''),h'',(h''',l''')\rangle^{\cap}llhhl'):N'' \bullet \\ \omega L\Sigma(l) = \{(\omega HI(h),\omega HI(h')),(\omega HI(h'),\omega HI(h))\} \land \\ \omega L\Sigma(l''') = \{(\omega HI(h''),\omega HI(h''')),(\omega HI(h'''),\omega HI(h'''))\} \land \\ \forall i,i+1:Nat \bullet \{i,i+1\}\subseteq inds llhhl \Rightarrow \\ let ((li,li'),hi,(hi'',li'')) = llhhl(i), \\ (\_,hj,(hj'',lj'')) = llhhl(i+1) in \\ \omega L\Omega(li) = \{(\omega HI(hi),\omega HI(hj))\} \land \\ \omega L\Omega(li') = \{(\omega HI(hj),\omega HI(hj))\} \land \\ \omega L\Omega(li') = \{(\omega HI(hj),\omega HI(hj))\} \land \\ \omega H\Omega(hi) = \{...\} ... 3 lines end
```

The last three lines of the axiom express that the pairs of intermidate links are one-way open in opposite directions and that hubs (...) are always open for through, on/off and reverse traffic.

■ End of Example

The determination facet has an own theory and anown development methodology.

4.2.4 Extension

By domain extension we understand an operation that applies to a (projected and possibly determined and instantiated) domain description, i.e., a (domain) requirements prescription, and yields a (domain) requirements prescription. The latter prescribes that a software system is to support, partially or fully, entities, operations, events and/or behaviours that we not feasible (or not computable in reasonable time) in a domain without computing support, but which are now are not only feasible but also computable in reasonable time.

Example 17. *Extension:* We extend the domain by introducing toll road entry and exit booths as well as electronic ticket hub sensors and actuators. There should now follow a careful narrative and formalisation of these three machines: the car driver/machine "dialogues" upon entry and exit as well as the sensor/car/actuator machine "dialogues" when cars enter hubs. The description should first, we suggest, be ideal; then it should take into account failures of booth equipment, failures of electronic tickets, failures of car drivers, and failures of sensors and actuators.

The extension facet has an own theory and anown development methodology.

4.2.5 Fitting

By domain requirements fitting we understand an operation which takes two or more (say n) domain requirements prescriptions, d_{r_i} , that are claimed to share entities, actions, events and/or behaviours and map these into n+1 domain requirements prescriptions, δ_{r_i} , where one of these, $\delta_{r_{n+1}}$ capture the shared phenomena and concepts and the other n prescriptions, δ_{r_i} , are like the n "input" domain requirements prescriptions, d_{r_i} , except that they now, instead of the "more-or-less" shared prescriptions, that are now consolidated in $\delta_{r_{n+1}}$, prescribe interfaces between δ_{r_i} and $\delta_{r_{n+1}}$ for $i:\{1..n\}$.

Example 18. Fitting: We assume three ongoing requirements development projects, all focused around road transport net software systems: (i) road maintenance, (ii) toll road car monitoring and (iii) bus services on ordinary plus toll road net. The main shared phenomenon is the road net, i.e., the links and the hubs. The consolidated, shared road net domain requirements prescription, $\delta_{r_{n+1}}$, is here suggested to become a prescription for a shared database for, abstractly speaking, two relations, one for hubs and one for links. Tuples of these relations then prescribe representation of all hub, respectively all link attributes – common to the three applications. Functions (including actions) on hubs and links become database queries and updates. Etc.

The *fitting* facet has an own theory and anown development methodology.

4.2.6 Discussion

The last page or so have very briefly surveyed and illustrated domain requirements. The reader should take cognizance of the fact that these are indeed "derived" from the domain description. They are not domain descriptions, but, once the business process re-engineering has been adopted and the required software has been installed, then the domain requirements become part of a revised domain description!

4.3 Interface Requirements

By interface requirements we understand such requirements which are concerned with the phenomena and concepts *shared* between the domain and the machine. Thus such requirements can only be expressed using terms from both

the domain and the machine. We tackle the problem of "deriving", i.e., constructing interface requirements by tackling four "smaller" problems: those of "deriving" interface requirements for entities, actions, events and behaviours.

4.3.1 Entity Interfaces

Entities that are shared between the domain and the machine must initially be input to the machine. Dynamically arising entities must likewise be input and all such machine entities must have their attributes updated, when need arise. Requirements for shared entities thus entail requirements for their representation and for their human/machine and/or machine/machine transfer dialogue.

Example 19. Shared Entities: Main shared entities are those of hubs and links. Above, in Example 18, it was suggested that some relational database be used for representing hubs in a relation and representing links in another relation. In an abstract representation of these relations, that is, in a first normal form representation, hub and link identifiers become the primary key. As for human input, some man/machine dialogue based around a set of visual display unit screens with fields for the input of hub, respectively link attributes can then be devised. Etc.

The shared entities facet has an own theory and anown development methodology.

4.3.2 Action Interfaces

By a shared action we mean an action that can only be partly computed by the machine. That is, the machine, in order to complete an action, may have to inquire with the domain (some measurable, time-varying entity attribute value, or some domain stake-holder) in order to proceed in its computation.

Example 20. Shared Actions: In order for a car driver to leave an exit toll both the following component actions must take place: the driver inserts the electronic pass in the exit toll booth machine. The machine scans and accepts the ticket and calculates the fee for the car journey from entry booth via the toll road net to the exit booth. The driver is alerted to the cost and is requested to pay this amount. Once paid the exit booth toll gate is raised. Notice that a number of details of the new support technology is left out. It could either be elaborated upon here, or be part of the system design.

■ End of Example

The *shared actions* facet has an own theory and anown development methodology.

4.3.3 Event Interfaces

By a shared event we mean an event whose occurrence in the domain need be communicated to the machine – and, vice-versa, an event whose occurrence in the machine need be communicated to the domain.

Example 21. Shared Events: The arrival of a car at a toll plaza entry booth is an event that must be communicated to the machine so that the entry booth may issue a proper pass (ticket). Similarly for the arrival at a toll plaza exit booth so that the machine may request the return of the pass and compute the fee. The end of that computation is an event that is communicated to the driver (in the domain) requesting that person to pay a certain fee. And so forth.

The *shared events* facet has an own theory and anown development methodology.

4.3.4 Behaviour Interfaces

By a shared behaviour we understand a sequence of zero, one or more shared actions and events.

Example 22. Shared Behaviour: A typical toll road net use behaviour is as follows: Entry at some toll plaza: receipt of electronic ticket, placement of ticket in special ticket "pocket" in front window, the raising of the entry booth toll gate; drive up to [first] toll road hub (with electronic registration of time of occurrence), drive down a selected link (with electronic registration of time of occurrence of entry to and exit from link), then a repeated number of zero, one or more toll road hub and link visits – some of which may be "repeats" – ending with a drive down from a toll road hub to a toll plaza with the return of the electronic ticket, etc. – cf. Example 21.

■ End of Example

The shared behaviour facet has an own theory and anown development methodology.

4.3.5 Discussion

The discussion of Sect. 4.2.6 carries over to this section. That is, once the machine has been installed it, the machine, is part of the new domain!

4.4 Machine Requirements

Since, basically, domains, other than the introspective machine domain itself, has no bearing on machine requirements we shall not cover this stage of requirements development other than saying that it consists of the following concerns: performance requirements (storage, speed, other resources), dependability requirements (availability, accessibility, integrity, reliability, safety, security), maintainability requirements (adaptive, extensional, corrective, perfective, preventive), portability requirements (development platform, execution platform, maintenance platform, demo platform) and documentation requirements. Only dependability seems to be subjectable to rigorous, formal treatment. We refer to [3, Vol. 3, Part V, Chap. 19, Sect. 19.6] for an extensive (30 page) survey.

The **discussions** of Sects. 4.2.6 and 4.3.5 carry over to this paragraph. That is, once the machine has been installed it, the machine, is part of the new domain!

5. DISCUSSION

[1] Domain Descriptions Are Not Normative: The description of, for example, "the" domain of the New York Stock Exchange would describe the set of rules and regulations governing the submission of sell offers and buy bids as well as those of clearing ('matching') sell offers and buy bids. These rules and regulations appears to be quite different from those of the Tokyo Stock Exchange [24]. A normative description of stock exchanges would abstract these rules so as to be rather un-informative. And, anyway, rules and regulations changes and business process re-engineering changes entities, actions, events and behaviours. For any given software development one may thus have to rewrite parts of existing domain descriptions, or construct an entirely new such description.

[2] "Requirements Always Change": This claim is often used as a hidden excuse for not doing a proper, professional job of requirements prescription, let alone "deriving" them, as we advocate, from domain descriptions. Instead we now make the following counterclaims [1] "domains are far more

stable than requirements" and [2] "requirements changes arise more as a result of business process re-engineering than as a result of changing stake-holder ideas". Cases (1), where it seems that domains are changing, are most often examples of business process re-engineering. Closer studies of a number of domain descriptions, for example of a financial service industry, reveals that the domain in terms of which an "ever expanding" variety of financial products are offered, are, in effect. based on a small set of very basic domain functions which have been offered for well-nigh centuries! We claim that thoroughly developed domain descriptions and thoroughly "derived" requirements prescriptions tend to stabilise the requirements re-design, but never alleviate it.

[3] What Have We Omitted: Our coverage of domain and requirements engineering has focused on modelling techniques for domain and requirements facets. We have omitted the important software engineering tasks of stake-holder identification and liaison, domain (and, to some extents also requirements) acquisition and analysis, terminologisation, and techniques for domain and requirements validation and verification. We refer, instead, to [3, Vol.3, Part IV (Chaps. 9, 12–14) and Part V (Chaps. 18, 20–23)].

[4] What Can Be Described and Prescribed: The issue of "what can be described" has been a constant challenge to philosophers. In [23] Russell covers his first Theory of Descriptions (stemming from the early 1900s), and in [22] a revision, as The Philosophy of Logical Atomism. The issue is not that straightforward. In [5, 6] we try to broach the topic from the point of view of the kind of domain engineering presented in this paper. Our approach is simple; perhaps too simple! We can describe what can be observed. We do so, first by postulating types of observable phenomena and of derived concepts; then by the introduction of observer functions and by axioms over these, that is, over values of postulated types and observers. To this we add defined functions; usually described by pre/post-conditions. The narratives refer to the "real" phenomena whereas the formalisations refer to related phenomenological concepts. The narrative/formalisation problem is that one can 'describe' phenomena without always knowing how to formalise them.

[5] What Have We Achieved – and What Not: Section 1.4 made some claims. We think we have substantiated them all, albeit ever so briefly.

Each of the domain facets (intrinsics, support technologies, management and organisation, rules and regulations, scrips [licenses and contracts] and human behaviour) and each of the requirements facets (projection, instantiation, determination, extension and fitting) provide rich grounds for both specification methodology studies and and for more theoretical studies [4].

[6] Relation to Other Works: The most obvious 'other' work is that of [19, M.A.Jackson: Problem Frames]. In [19] Jackson, like is done here, departs radically from conventional requirements engineering. In his approach understandings of the domain, the requirements and possible software designs are arrived at, not hierarchically, but in parallel, interacting streams of decomposition. Thus the 'Problem

Frame' development approach iterates between concerns of domains, requirements and software design. "Ideally" our approach pursues domain engineering prior to requirements engineering, and, the latter, prior to software design. But see next

[7] "Ideal" Versus Real Developments: The term 'ideal' has been used in connection with 'ideal development' from domain to requirements. We now discuss that usage. Ideally software development could proceed from developing domain descriptions via "deriving" requirements prescriptions to software design, each phase involving extensive formal specifications, verifications (formal testing, model checking and theorem proving) and validation. More realistically less comprehensive domain description development (D) may alternate with both requirements development (R) work and with software design (S) – in some hopefully controlled, contained "spiralling" manner and such that it is at all times clear which development step is what: \mathcal{D} , \mathcal{R} or \mathcal{S} !

[8] $\mathcal{D}, \mathcal{S} \models \mathcal{R}$: This reads: in a proof of correctness of \mathcal{S} oftware design with respect to \mathcal{R} equirements prescriptions one often has to refer to assumptions about the \mathcal{D} omain. Formalising our understandings of the \mathcal{D} omain, the \mathcal{R} equirements and the \mathcal{S} oftware design enables proofs that the software is rught and the \mathcal{D} omain formalisation of the "derivation" of \mathcal{R} equirements from \mathcal{D} omain specifications help ensure that it is the right software [10].

[9] Domain Versus Ontology Engineering: In the information science community an ontology is a "formal, explicit specification of a shared conceptualisation". Most of the information science ontology work seems aimed primarily at axiomatisations of properties of entities. Apart from that there are many issues of "ontological engineering" that are similar to "our kind" of domain engineering; but then, we claim, that domain engineering goes well beyond ontological engineering and makes free use of whatever formal specification languages are needed, cf. Sect. 1.3.

6. CONCLUSION

We have put forward the methodological steps of another approach to requirements engineering than currently 'en vogue'. We claim that our approach, as it follows from the **dogma** expressed in Sect. 1, is logical, that is, follows as a necessity. The 'triptych' approach has been in partial use since the early 1990s, notably at the United Nations University's International Institute for Software Technology (www.iist.-unu.edu). This paper presents this triptych is a clear form than presented in [3]. The (six) domain, the (five) domain requirements and the (4) interface requirements facets each have nice theories and each has a simple set of methodological principles and techniques – some covered in [3, 4, 5, 6] others to be further researched.

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7. BIBLIOGRAPHICAL NOTES

Section 1.3 gives most relevant references to formal specification languages (techniques and tools) that cover the

spectrum of domain and requirements specification, refinement and verification. The recent book on Logics of Specification Languages [7] covers ASM, B/event B, CafeObj, CASL, DC, RAISE, TLA+, VDM and Z.

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