

# Weekplan: Hashing

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## References and Reading

[1] Notes from Aarhus, Peter Bro Miltersen.

[2] Scribe notes from MIT.

[3] Universal Classes of Hash Functions, J. Carter and M. Wegman, J. Comp. Sys. Sci., 1977.

[4] Storing a Sparse Table with  $O(1)$  Worst Case Access Time, M. Fredman, J. Komlos and E. Szemerédi, J. ACM., 1984.

[5] Notes on Discrete Probability, Jeff Erickson.

We recommend reading [1] and [2] in detail. [3] and [4] provide background on universal and perfect hashing. [5] provides a concise refresh of basic discrete probability.

## Exercises

**1 [w] Streaming Statistics** An IT-security friend of yours wants a high-speed algorithm to count the number of *distinct* incoming IP-addresses in his router to help detect denial of service attacks. Can you help him?

**2 [w] Dense Set Hashing** A set  $S \subseteq U = \{0, \dots, u-1\}$  is called *dense* if  $|S| = \Theta(u)$ . Suggest a simple and efficient dictionary data structure for dense sets.

**3 [w] Multi-Set Hashing** A multi-set is a set  $M$ , where each element may occur multiple times. Design an efficient data structure supporting the following operations:

- $\text{add}(x)$ : Add an(other) occurrence of  $x$  to  $M$ .
- $\text{remove}(x)$ : Remove an occurrence of  $x$  from  $M$ . If  $x$  does not occur in  $M$  do nothing.
- $\text{report}(x)$ : Return the number of occurrences of  $x$ .

**4 Properties of Universal Hashing** Let  $h \in H$  be a hash function from a universal family mapping  $U = \{0, \dots, u-1\}$  to  $M = \{0, \dots, m-1\}$ . Solve the following exercises.

**4.1** If  $h$  has no collision on  $U$ , how large must  $m$  be?

**4.2** Suppose  $m \geq u$ . Is the identity function  $f(x) = x$  a universal hash function?

**4.3** A family  $G$  of hash functions mapping  $U$  to  $M$  is *family of pair-wise independent hash function* if for any  $g \in G$ ,

$$\Pr(g(x) = \alpha \wedge g(y) = \beta) = 1/m^2 \quad \forall x \neq y \in U, \quad \forall \alpha, \beta \in M.$$

Show that any family of pairwise independent hash functions is a family of universal hash functions.

**5 Linear Space Hashing** The chained hashing solution for the dynamic dictionary problem presented assume that  $m = \Theta(n)$ . Solve the following exercises.

- 5.1 What is the space and time of chained hashing without this assumption? State your answer in terms of  $n$  and  $m$ .
- 5.2 Suppose that we do not know  $n$  in advance (as in the exercise streaming statistics where we do not know how many distinct IP-address we will see). Give a solution that achieves  $O(n)$  space and constant time without assuming  $m = \Theta(n)$ . *Hint*: Think dynamic arrays.

**6 Graph Adjacency** Let  $G$  be a graph with  $n$  vertices and  $m$  edges. We want to represent  $G$  efficiently and support the following operation.

- $\text{adjacent}(v, w)$ : Return true if nodes  $v$  and  $w$  are adjacent and false otherwise.

Solve the following exercises:

- 6.1 Analyse the space and query time in terms of  $n$  and  $m$  for the classic adjacency matrix and adjacency list representation.
- 6.2 Design a data structure that improves both the adjacency matrix and adjacency list.

**7 Perfect Hashing Analysis** Consider the 2-level FKS perfect hashing scheme. A friend suggest the following two "optimizations" to the data structure. What happens to the performance of the data structure for each of these?

- 7.1 Modify level 1 of the data structure to map  $U$  to an array of size  $n\sqrt{n}$  instead of  $n$  to further decrease the probability of collisions.
- 7.2 Replace the universal hash function with a faster *near-universal hash function* on both levels. Near-universal hashing is the same as universal hashing except that  $\leq 1/m$  guarantee on the probability is changed to  $\leq 2/m$ .

**8 Lost Integer Puzzles** Suppose that you receive a stream of  $n-1$  distinct integers from the set  $\{1, \dots, n\}$ , i.e., the stream consists of all of  $\{1, \dots, n\}$  except a single missing integer. We want a space-efficient algorithm that efficiently computes this integer during a single pass over the input stream. Solve the following exercises:

- 8.1 [ $w$ ] Give a simple data structure to find the lost integer using  $O(n)$  space.
- 8.2 [ $*$ ] Show how to find the lost integer using  $O(1)$  space.
- 8.3 [ $**$ ] Suppose there are now two lost integers. Show how to find them using  $O(1)$  space.

**9 Basic Probability Theory Refresh Bonus** In case your knowledge of probability theory is rusty. Solve the following self-help exercises.

- 9.1 Prove linearity of expectation.
- 9.2 Prove that the expectation of the *indicator function* for  $h(x) = h(y)$  (1 if  $h(x) = h(y)$  and 0 otherwise) is equal to the probability that  $h(x) = h(y)$ .
- 9.3 Show that the expected number of trials to get a perfect hashing function using an array of size  $n^2$  is  $\leq 2$ .